Approximating Shortest Connected Graph Transformation for Trees

SOFSEM 2020

Nicolas Bousquet, Alice Joffard

January 21, 2020

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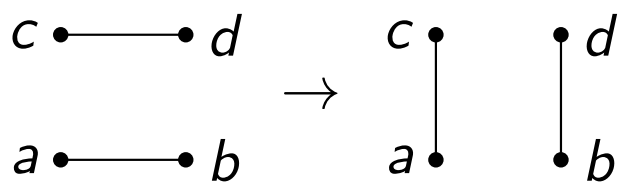

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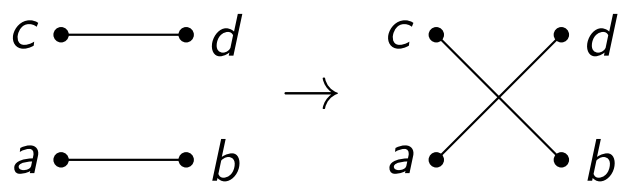

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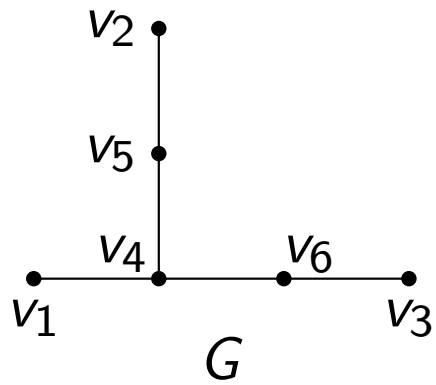

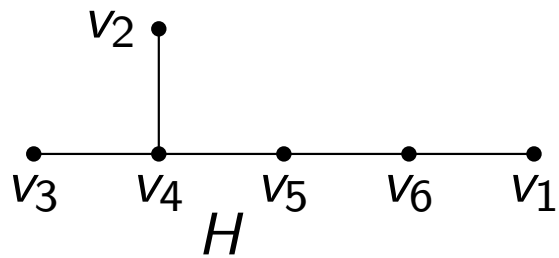

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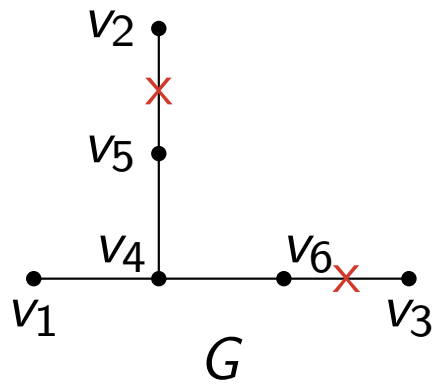

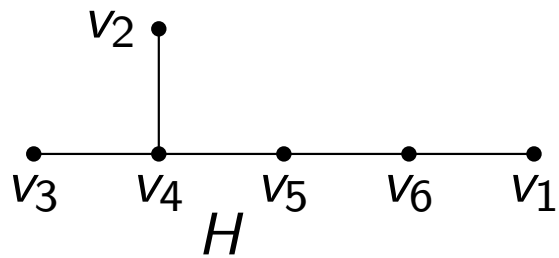

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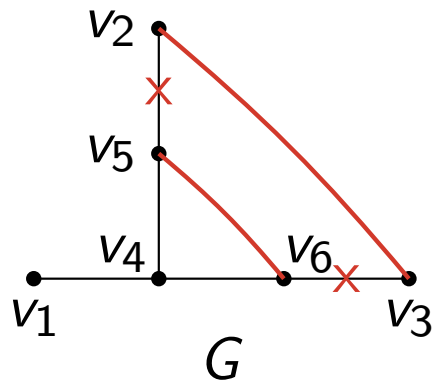

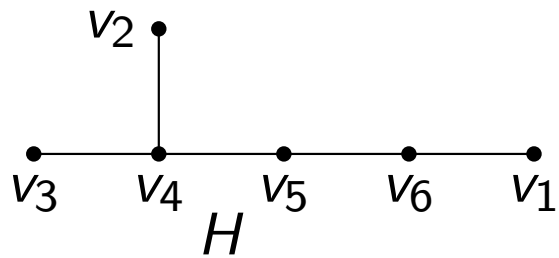

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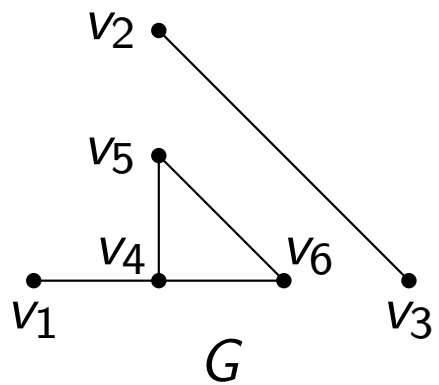
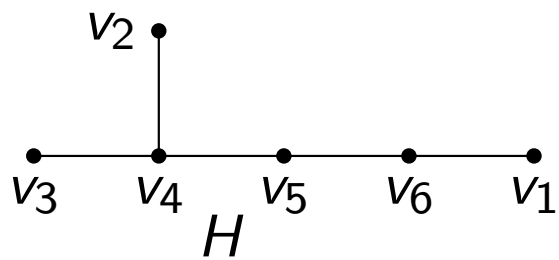


G

DISCONNECTED



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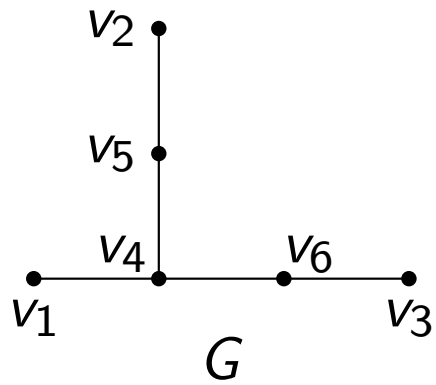

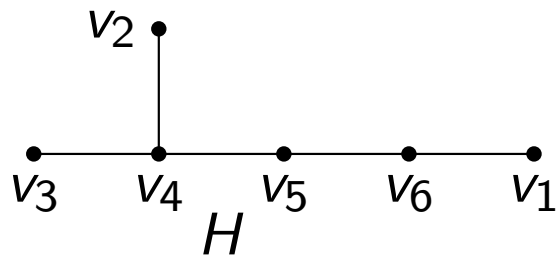

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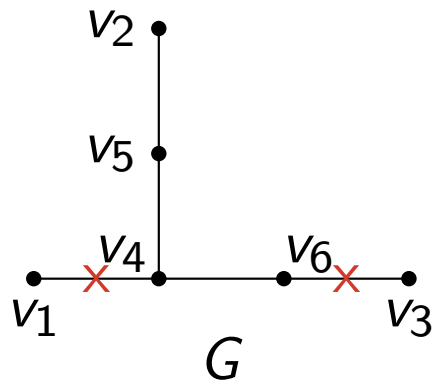

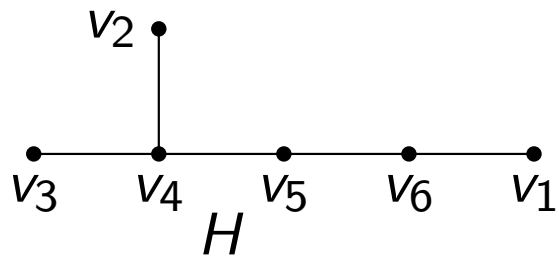

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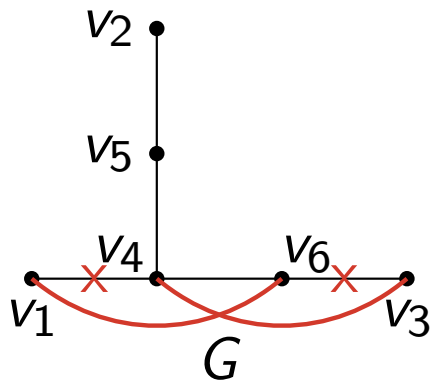

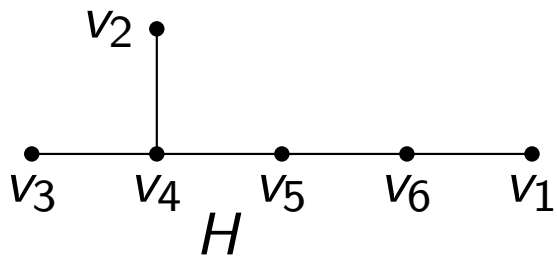

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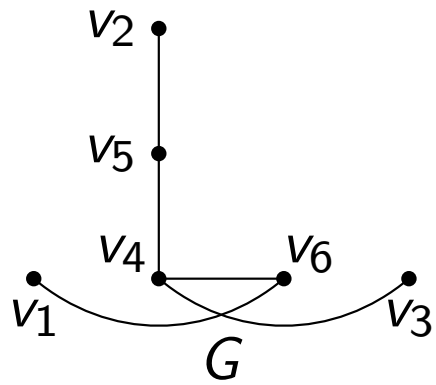

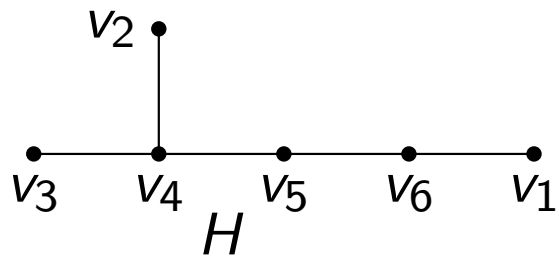

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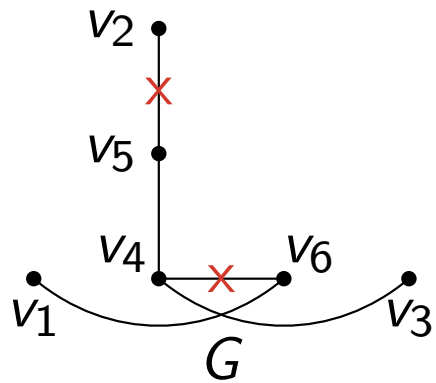

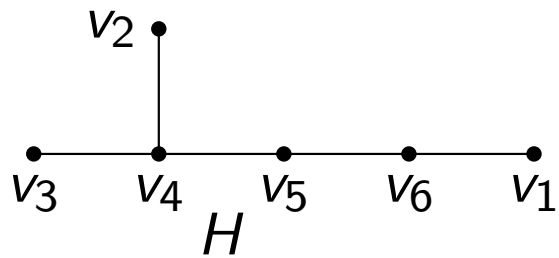

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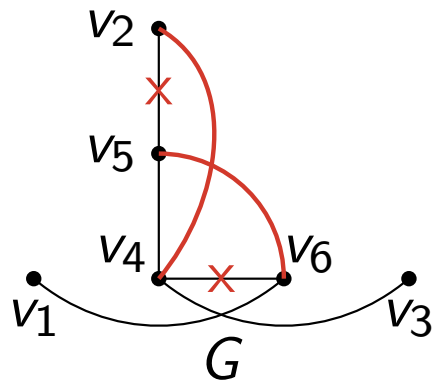

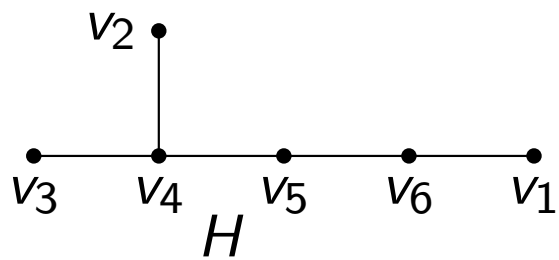

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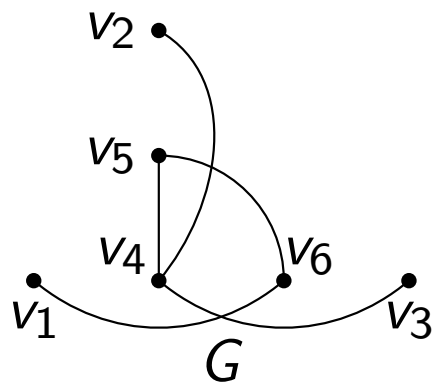

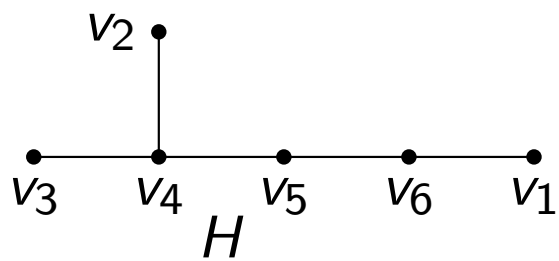

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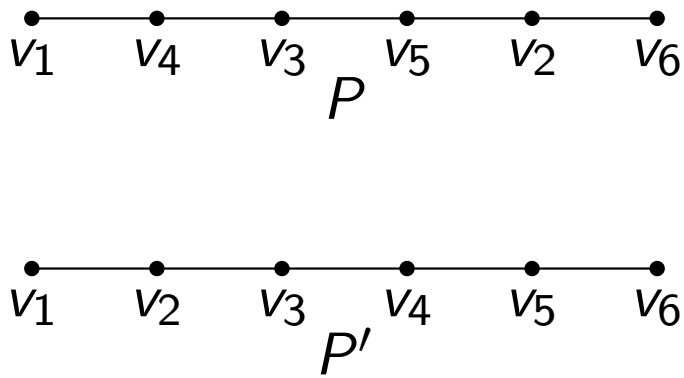

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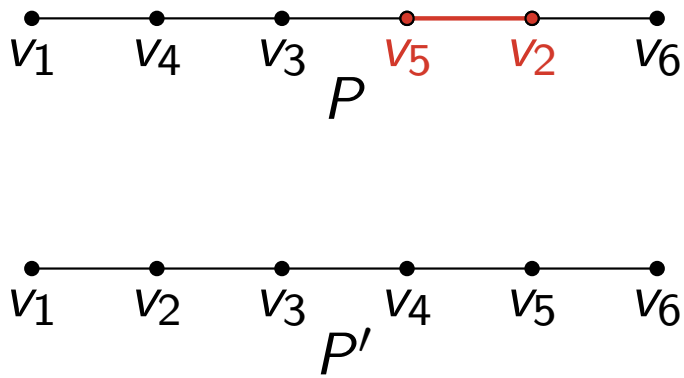

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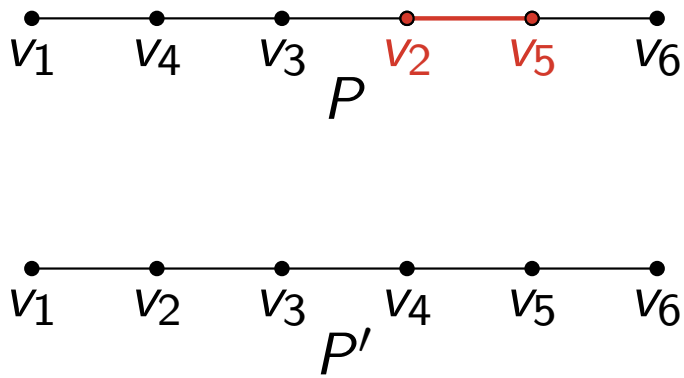

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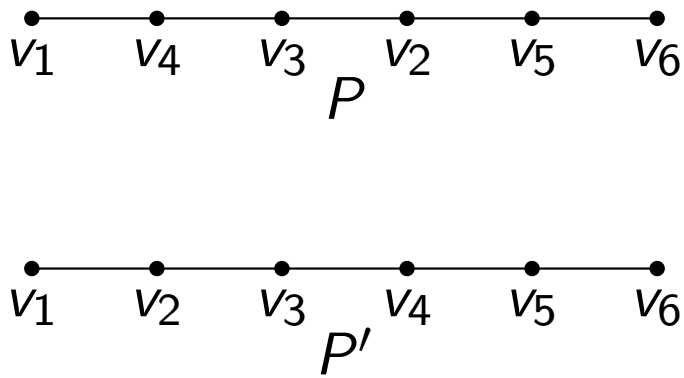

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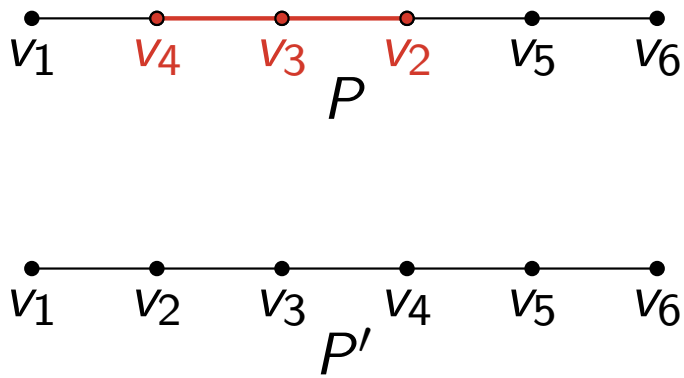

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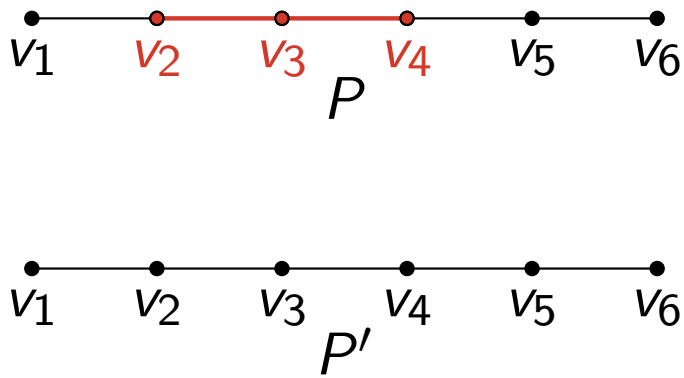

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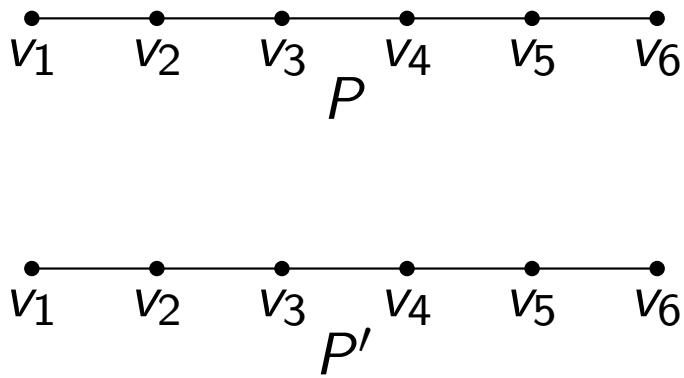

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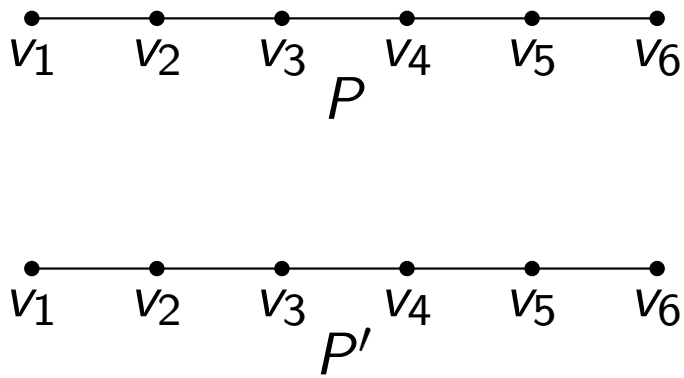

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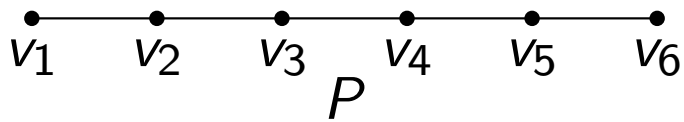

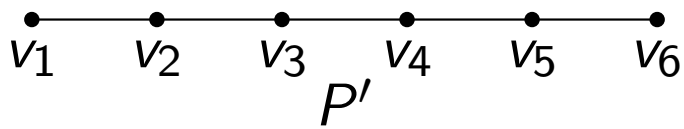

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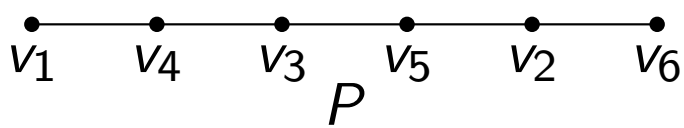

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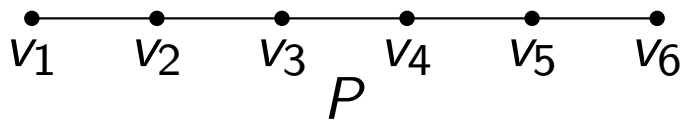

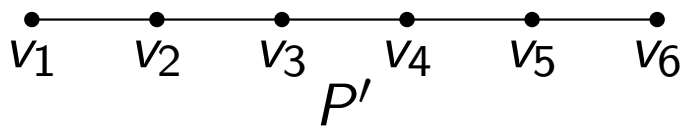

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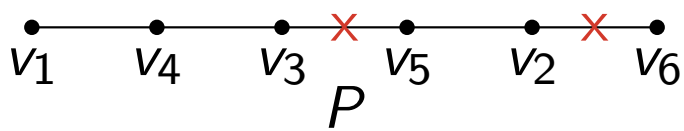

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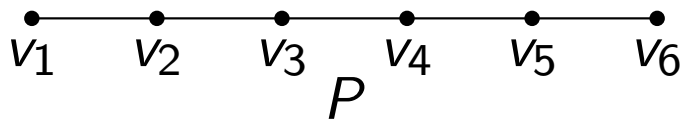

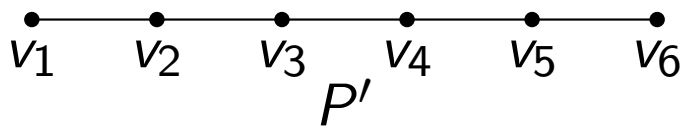

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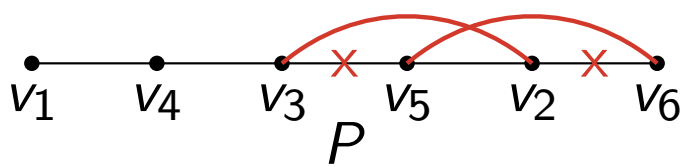

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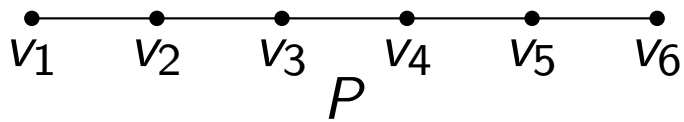

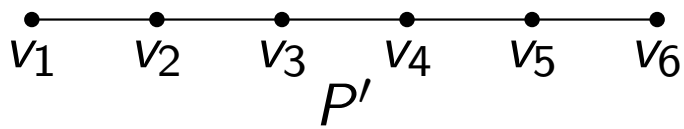

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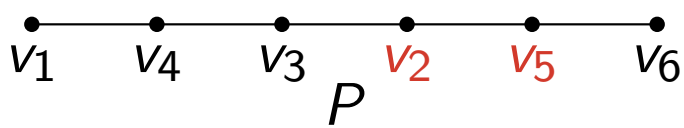

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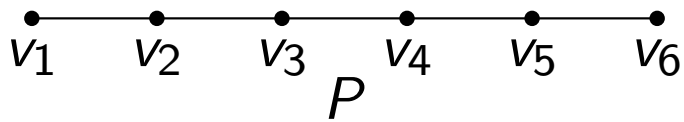

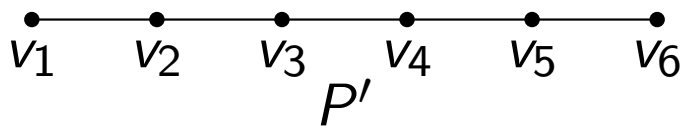

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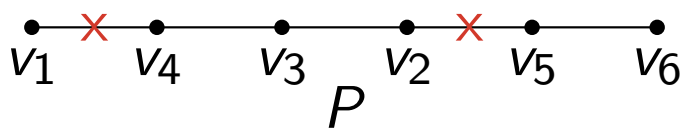

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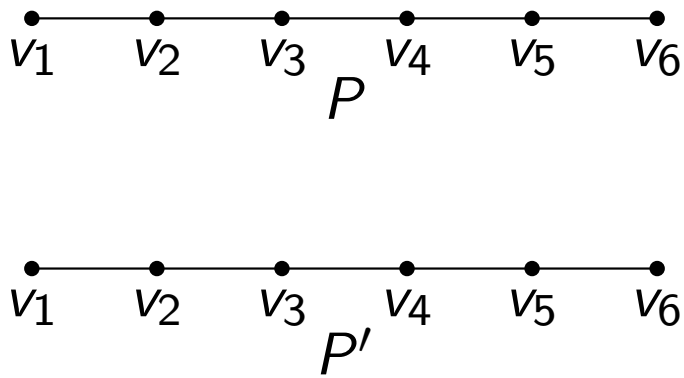

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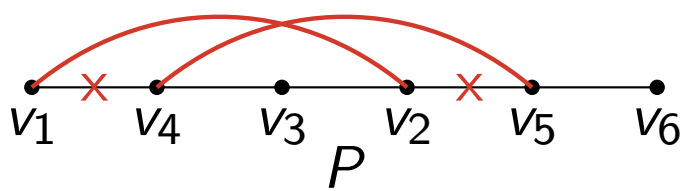

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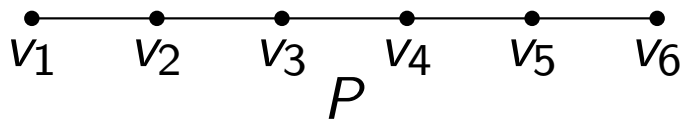

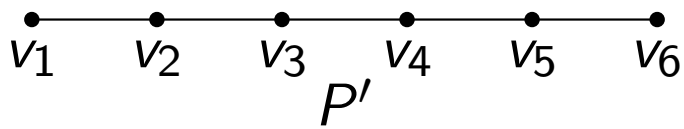

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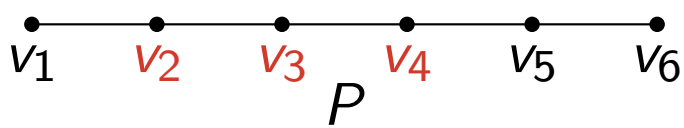

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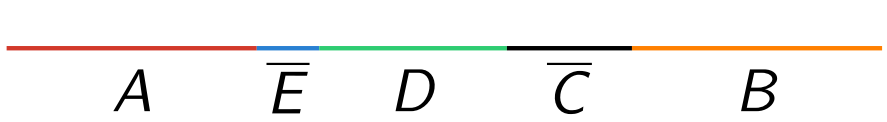

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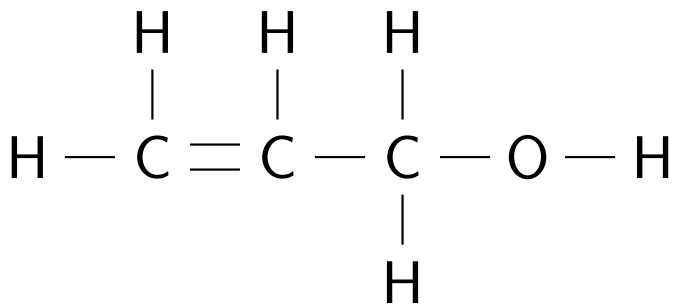
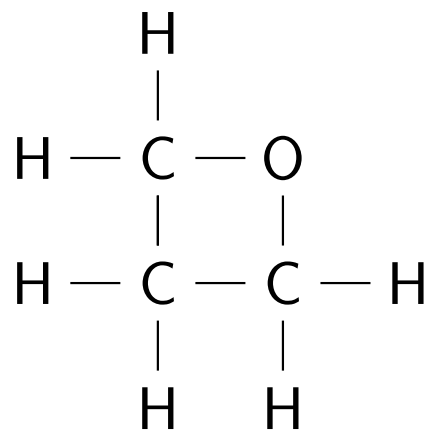


allyl alcohol C_3H_6O



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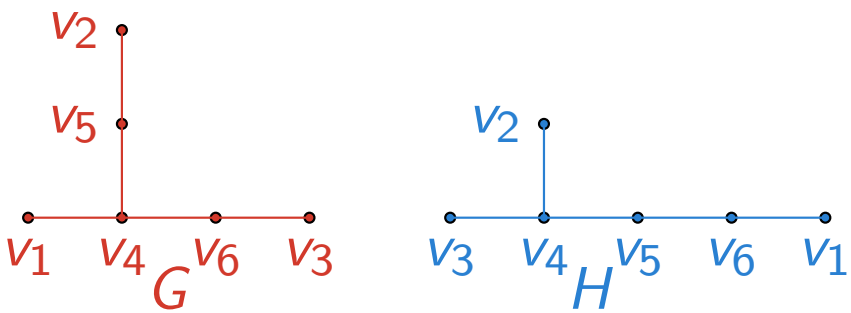

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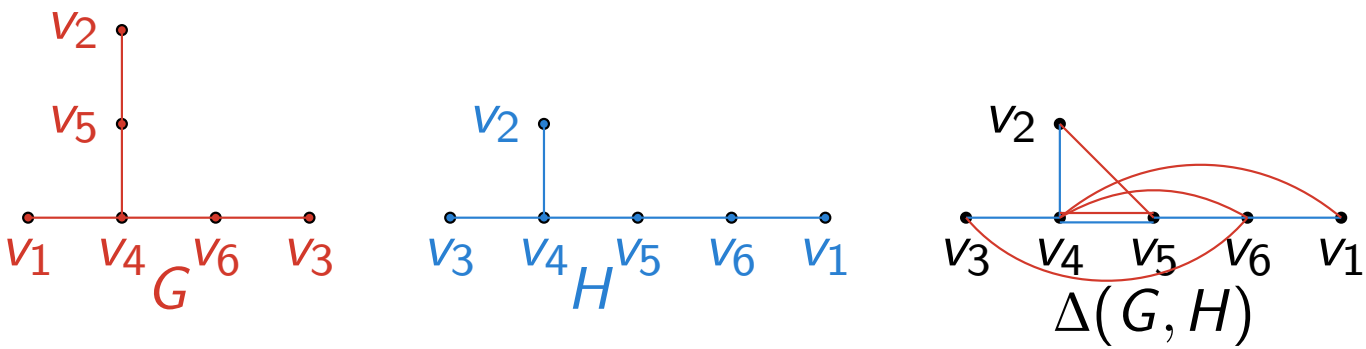

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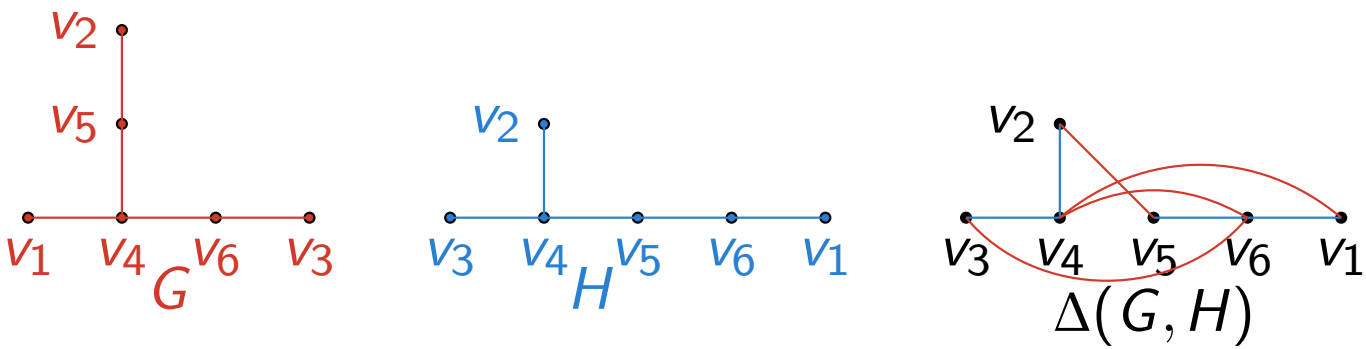

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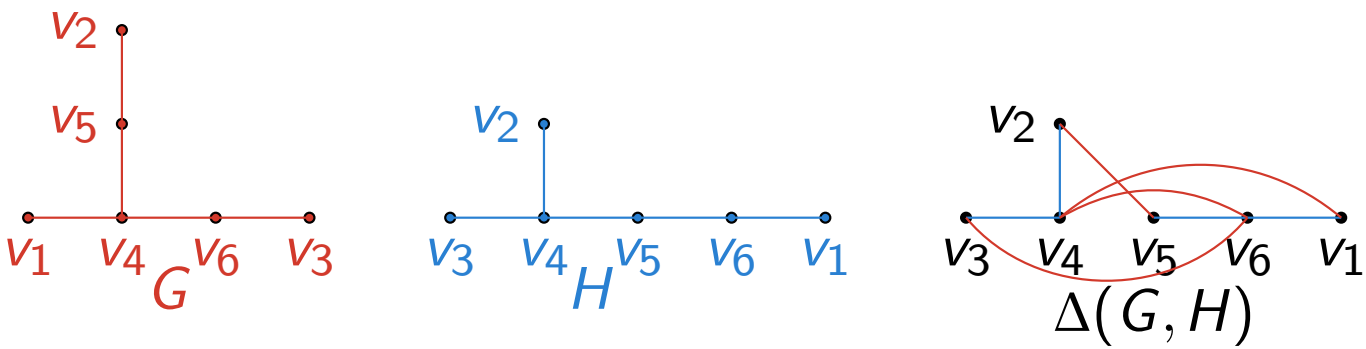

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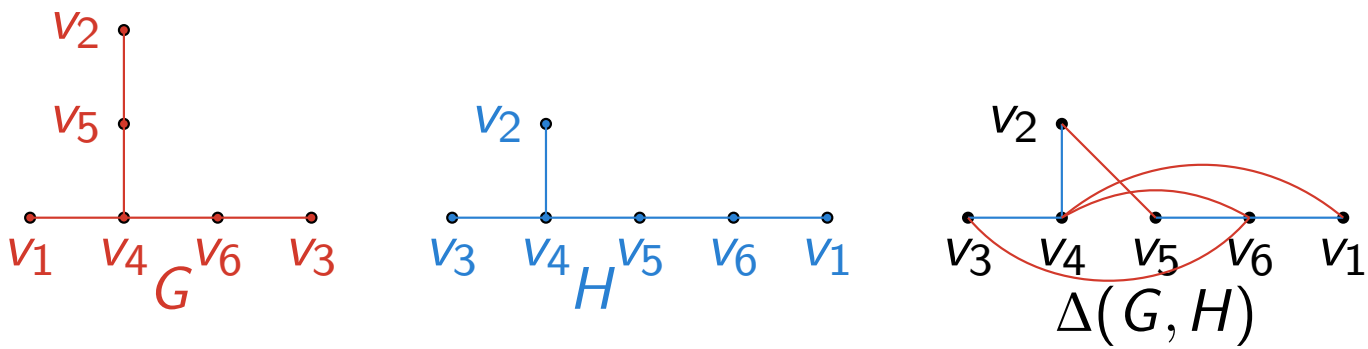

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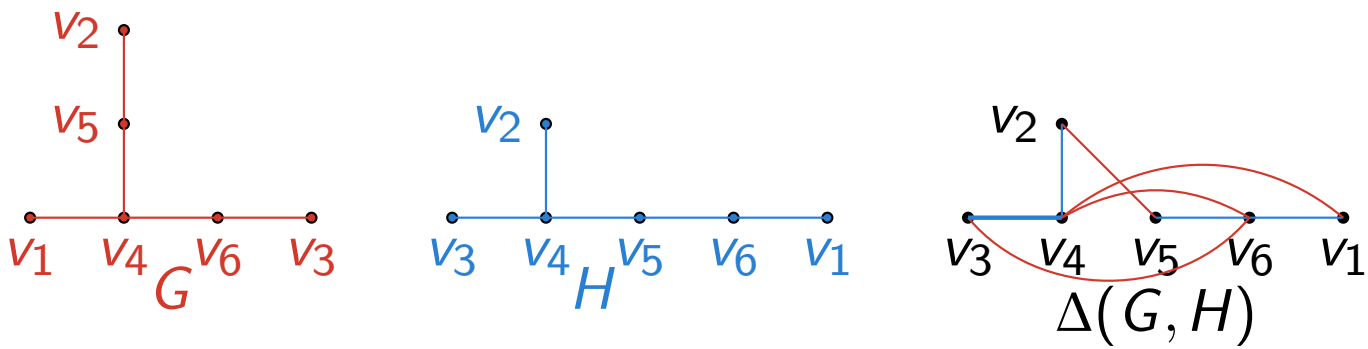

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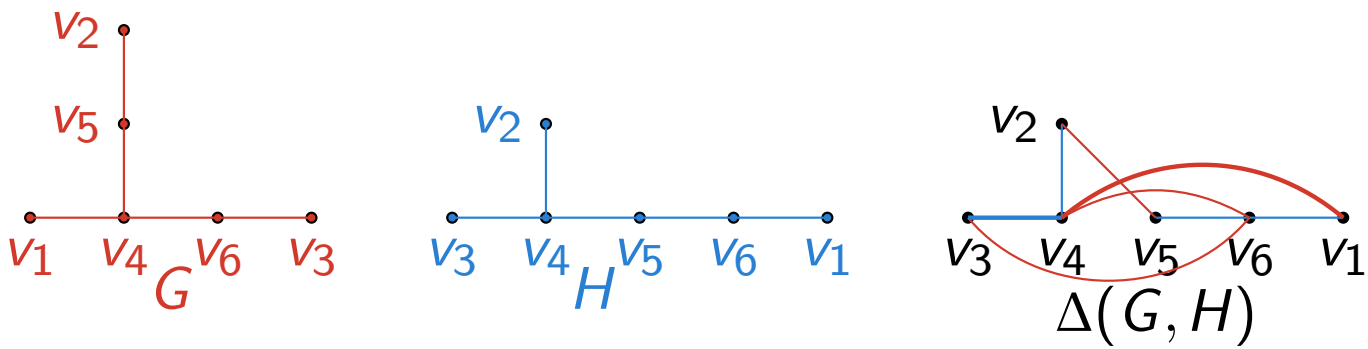

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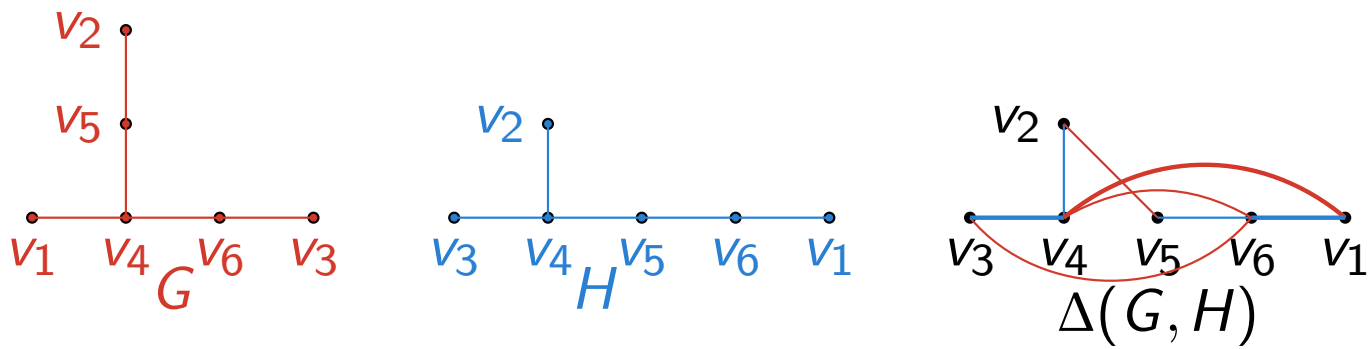

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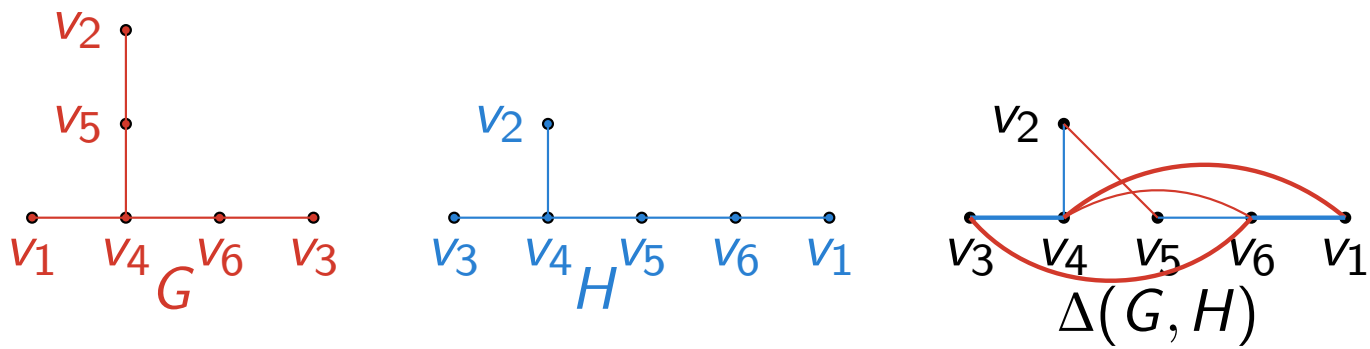

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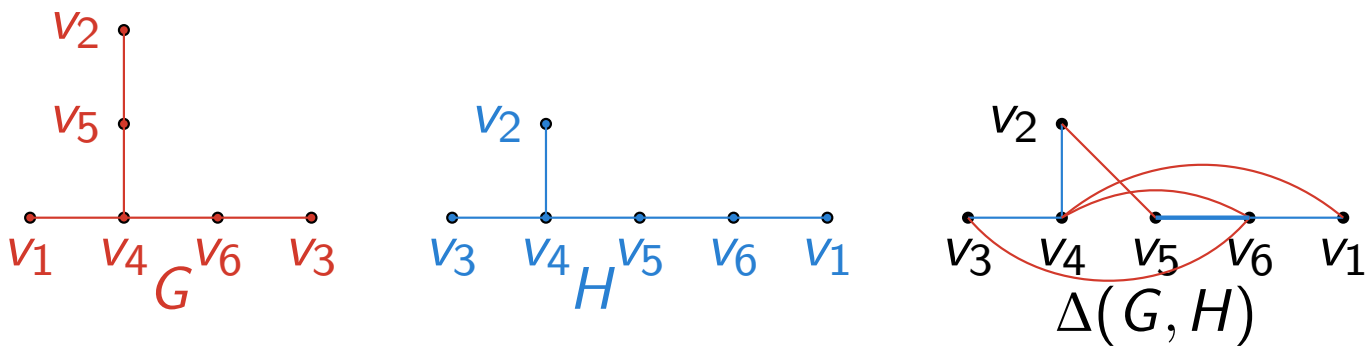

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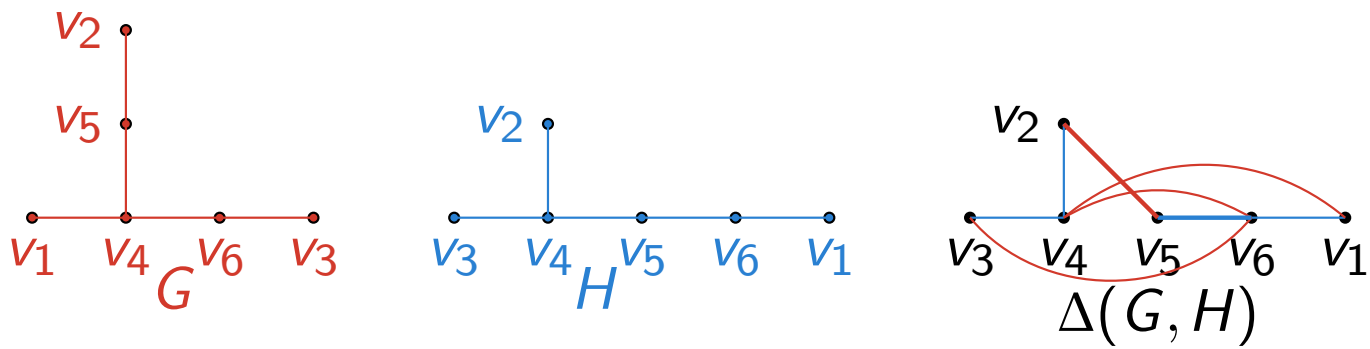

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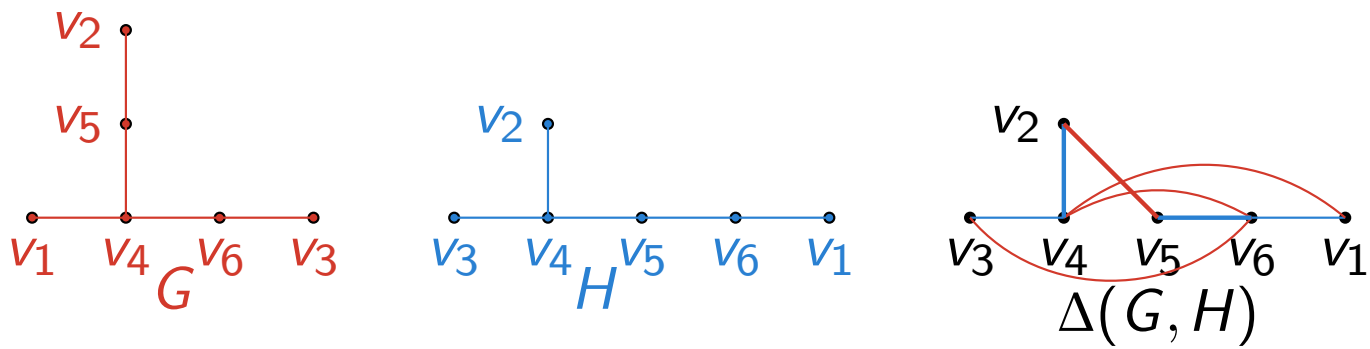

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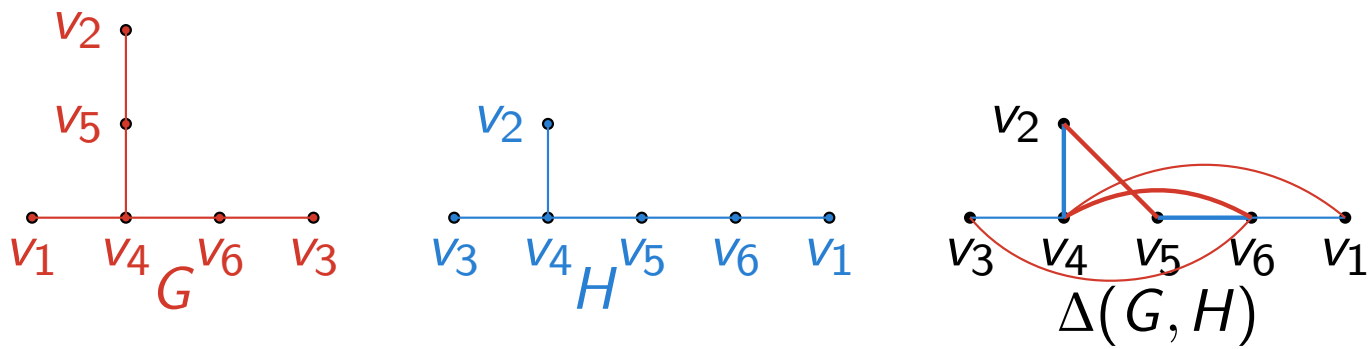

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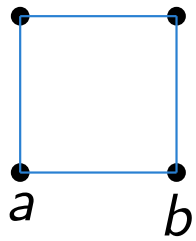

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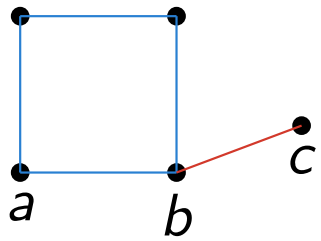

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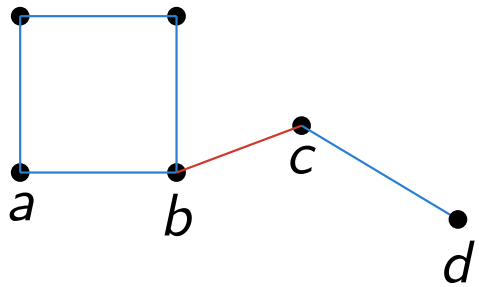

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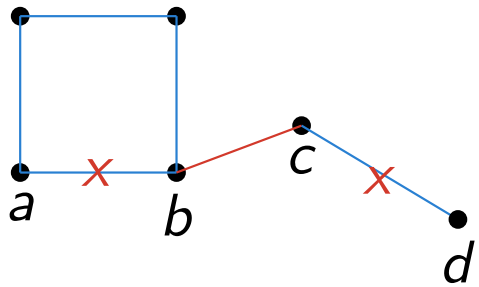

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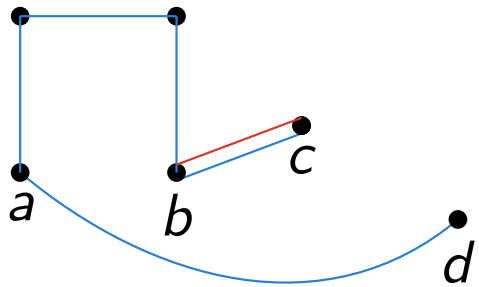

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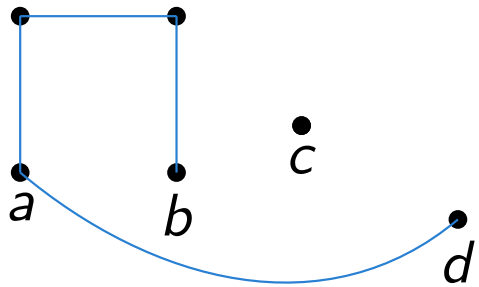

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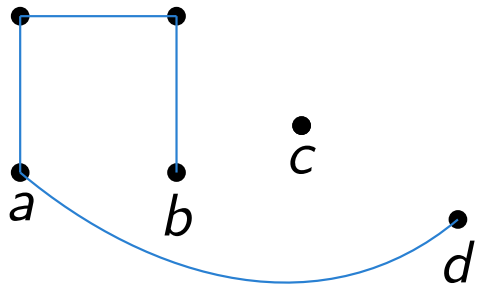

- Consequences of a flip on a tree:

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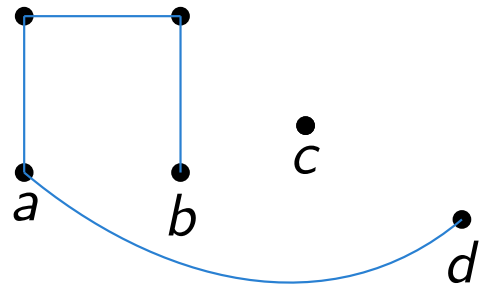

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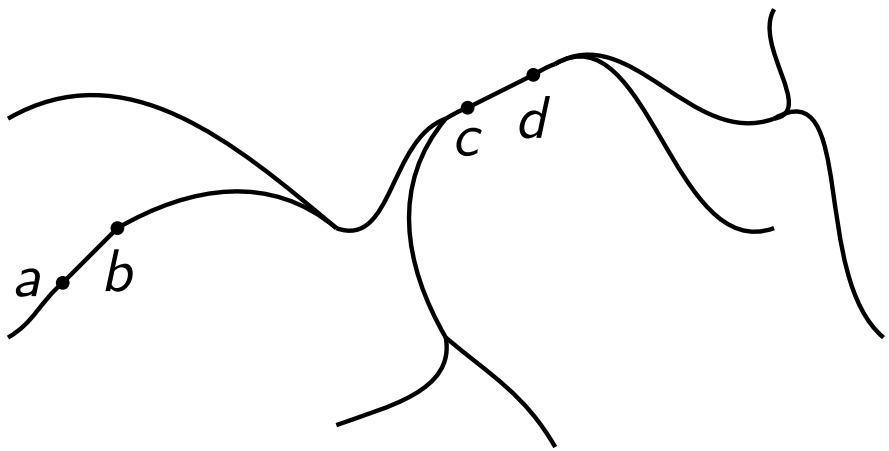

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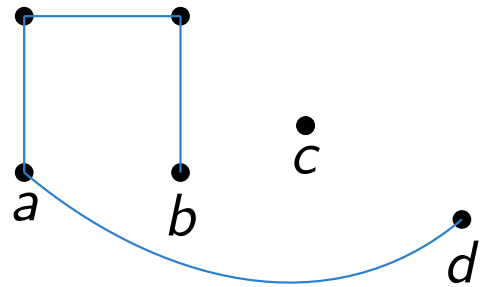

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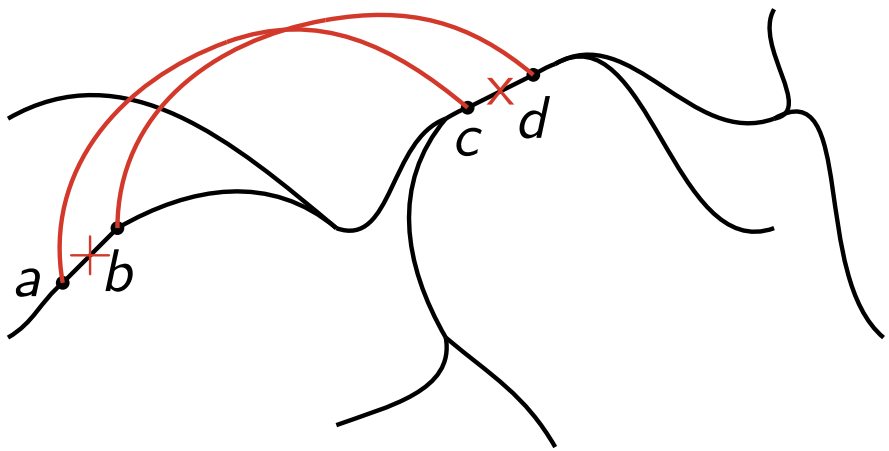

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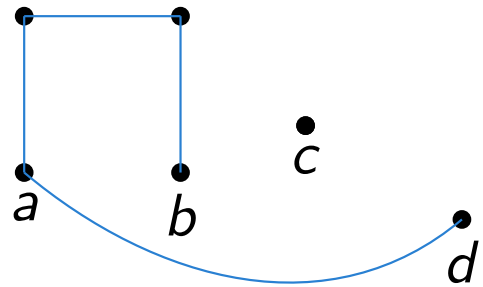

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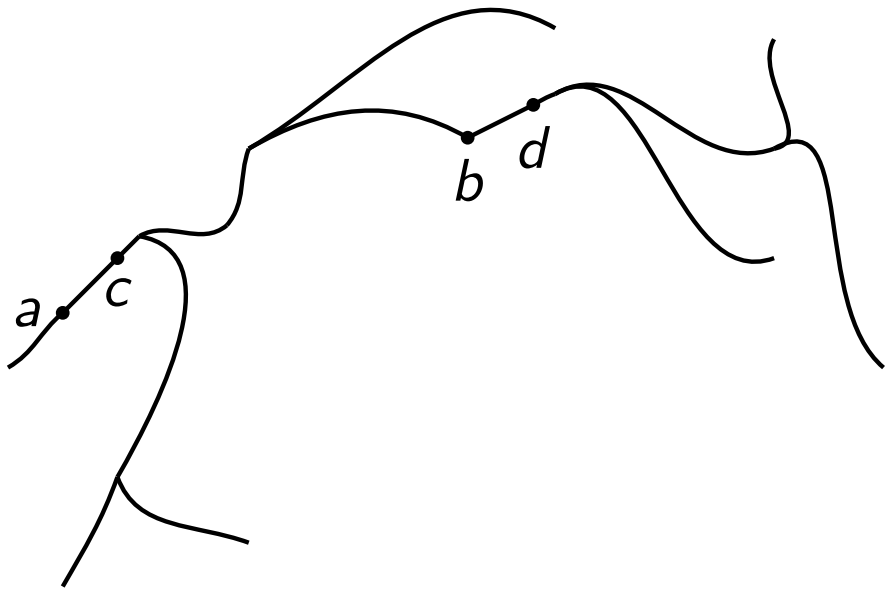

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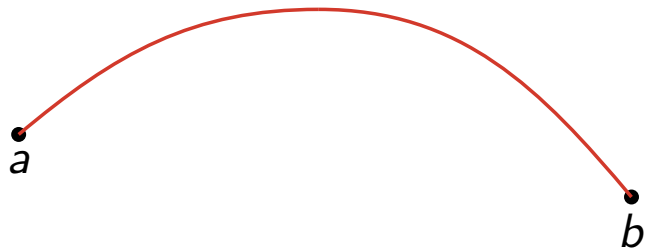

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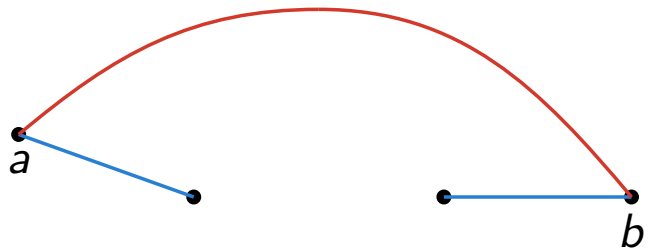

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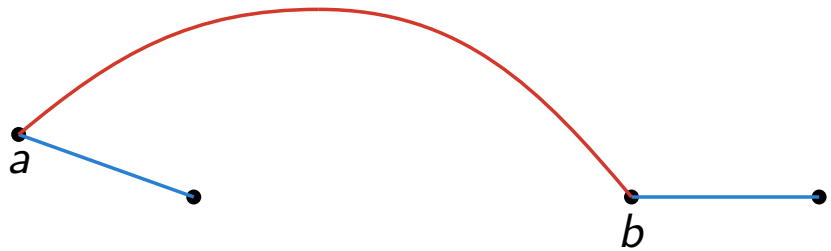

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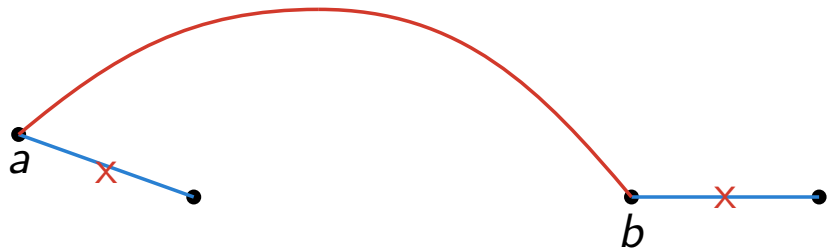

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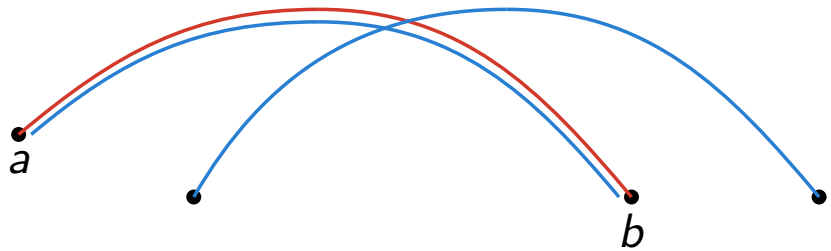

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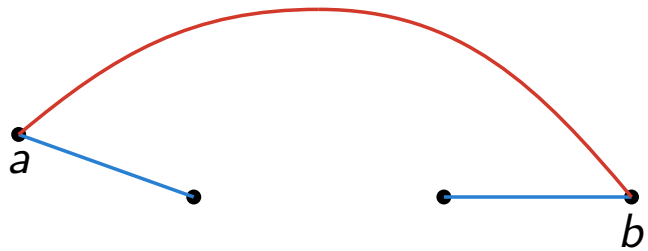

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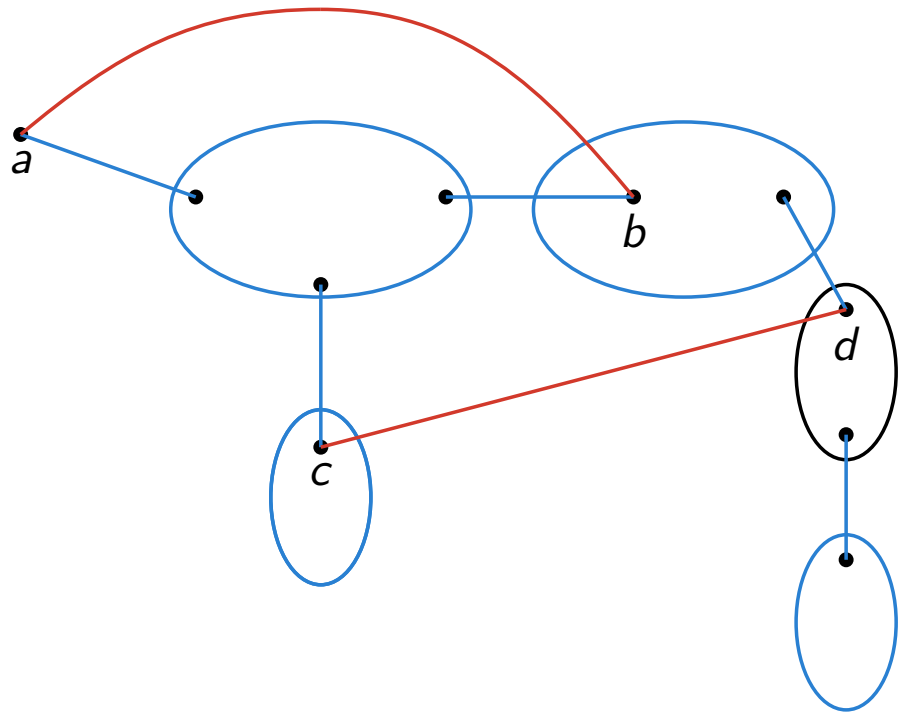

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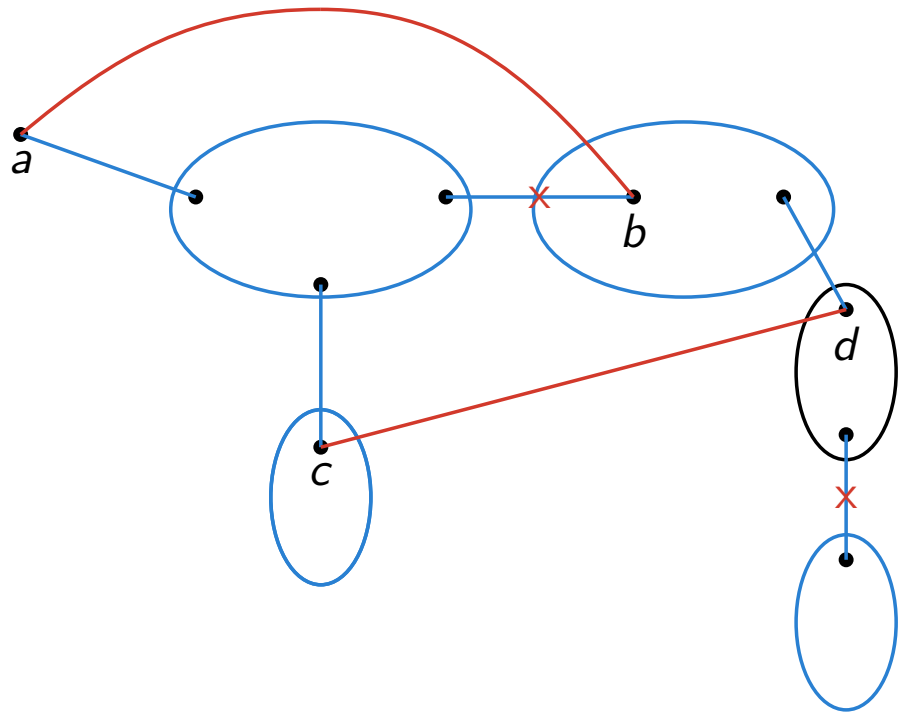

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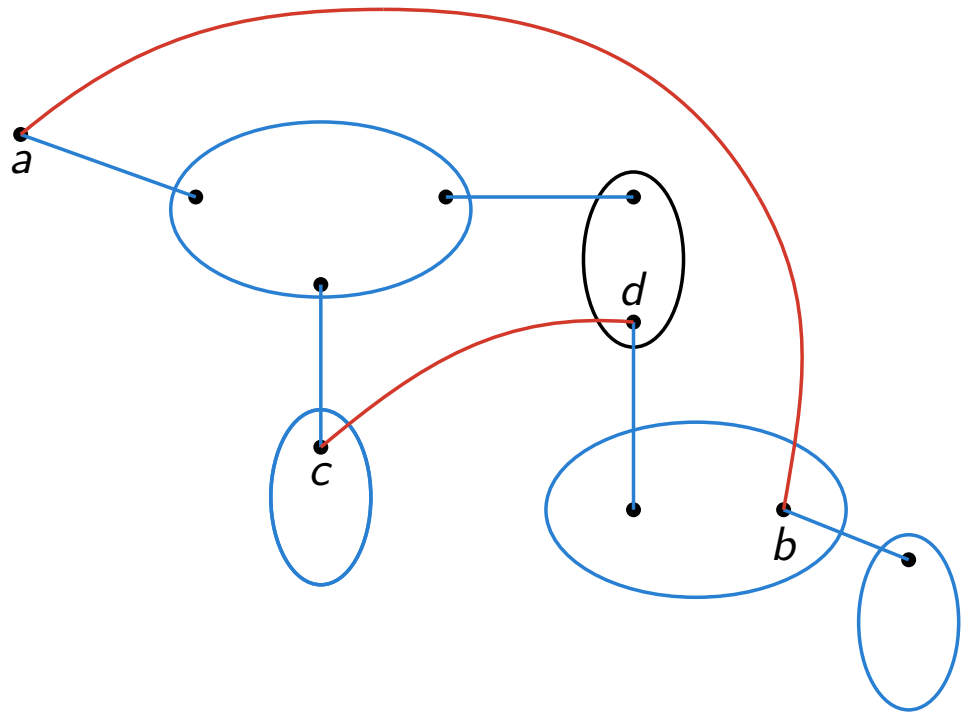

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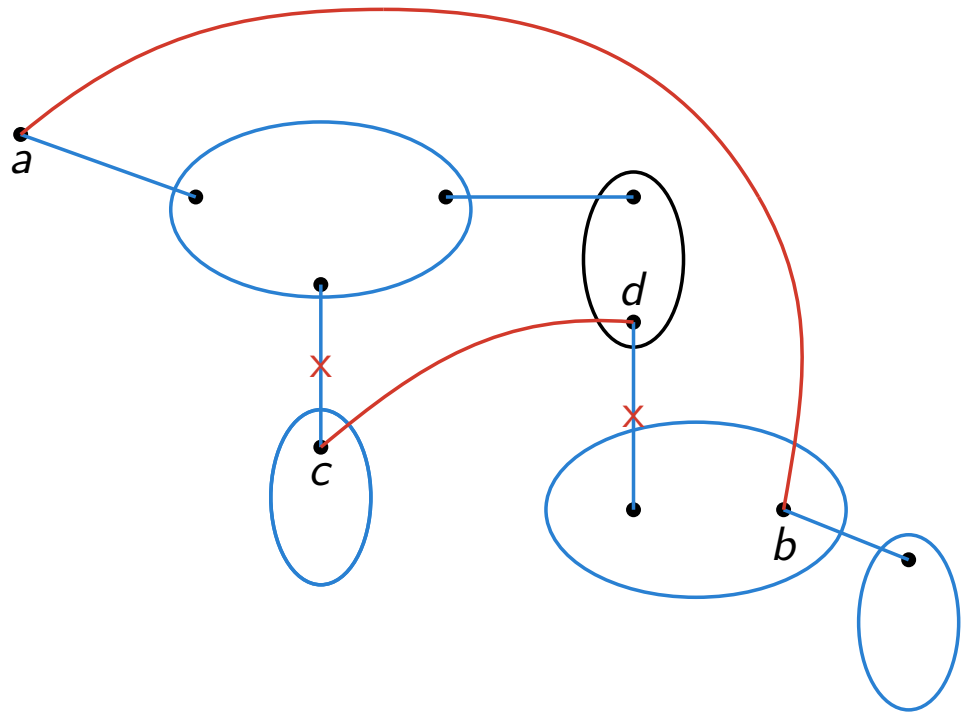

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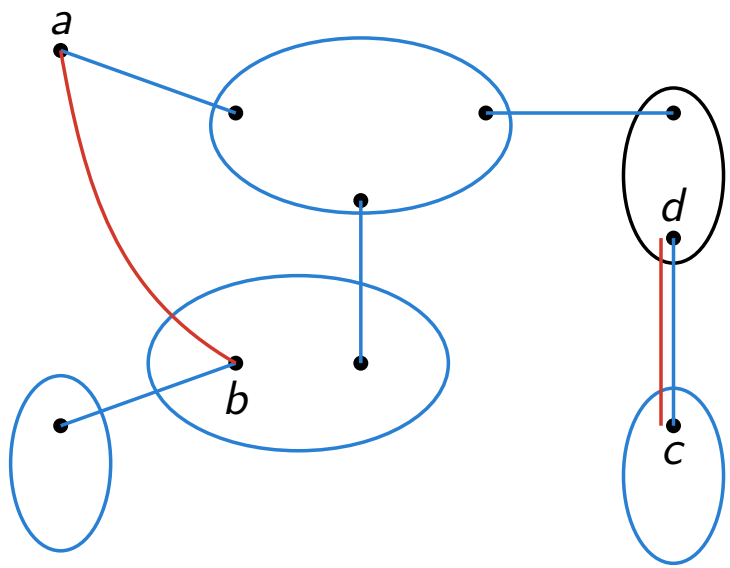

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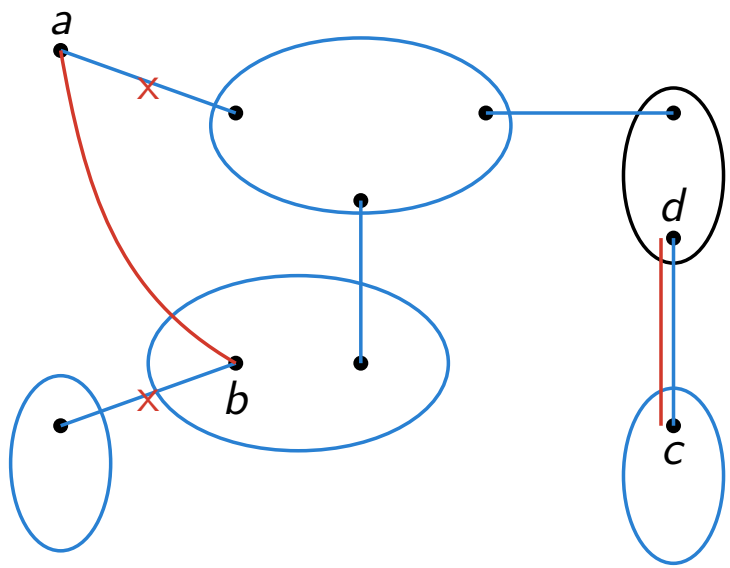

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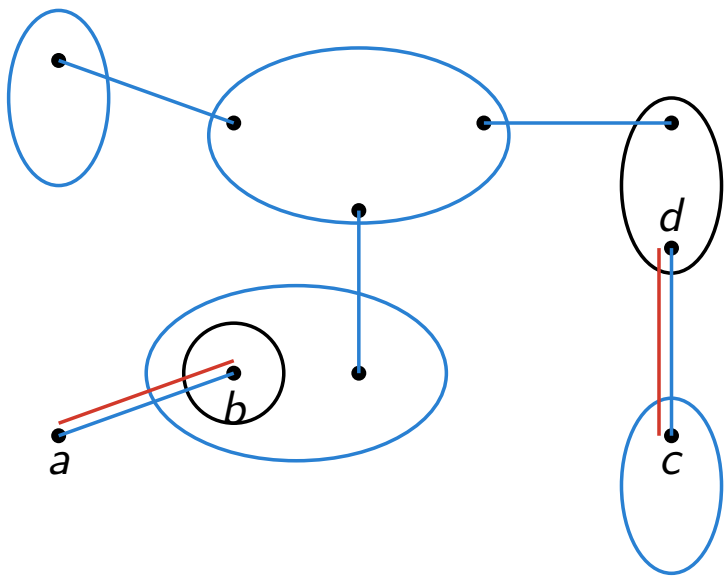

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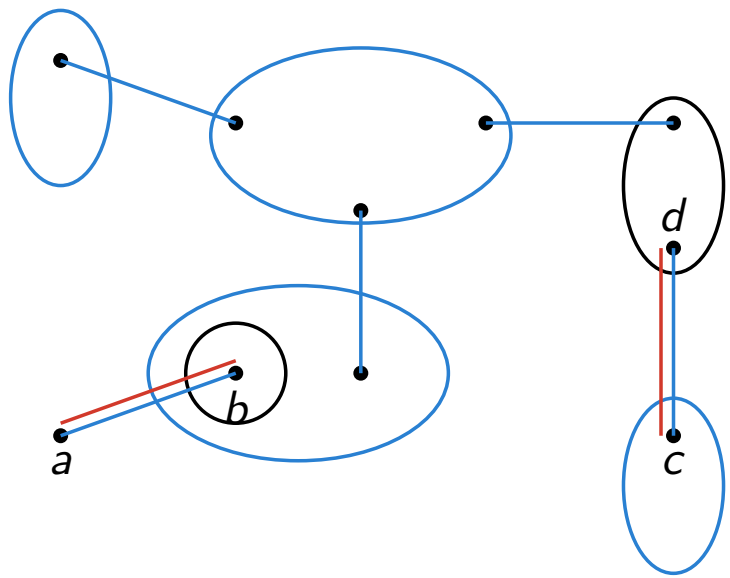

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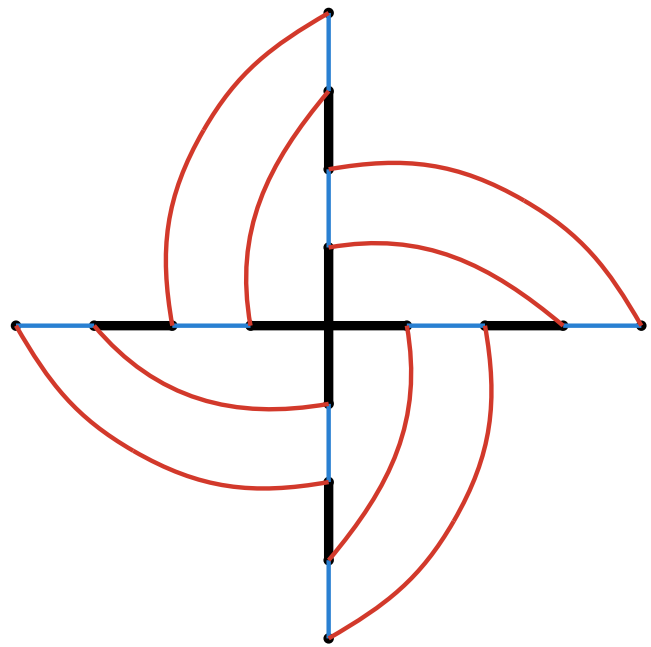

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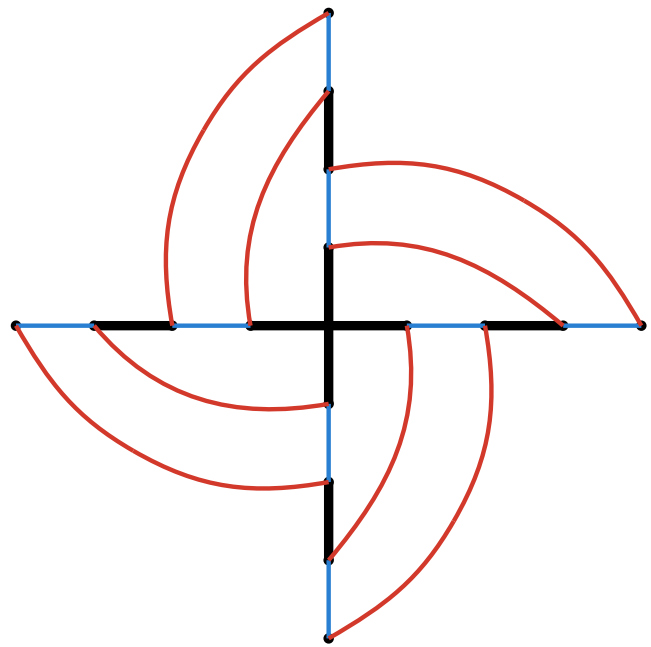

- With this lower bound, can not do better than a 1.5-approximation

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