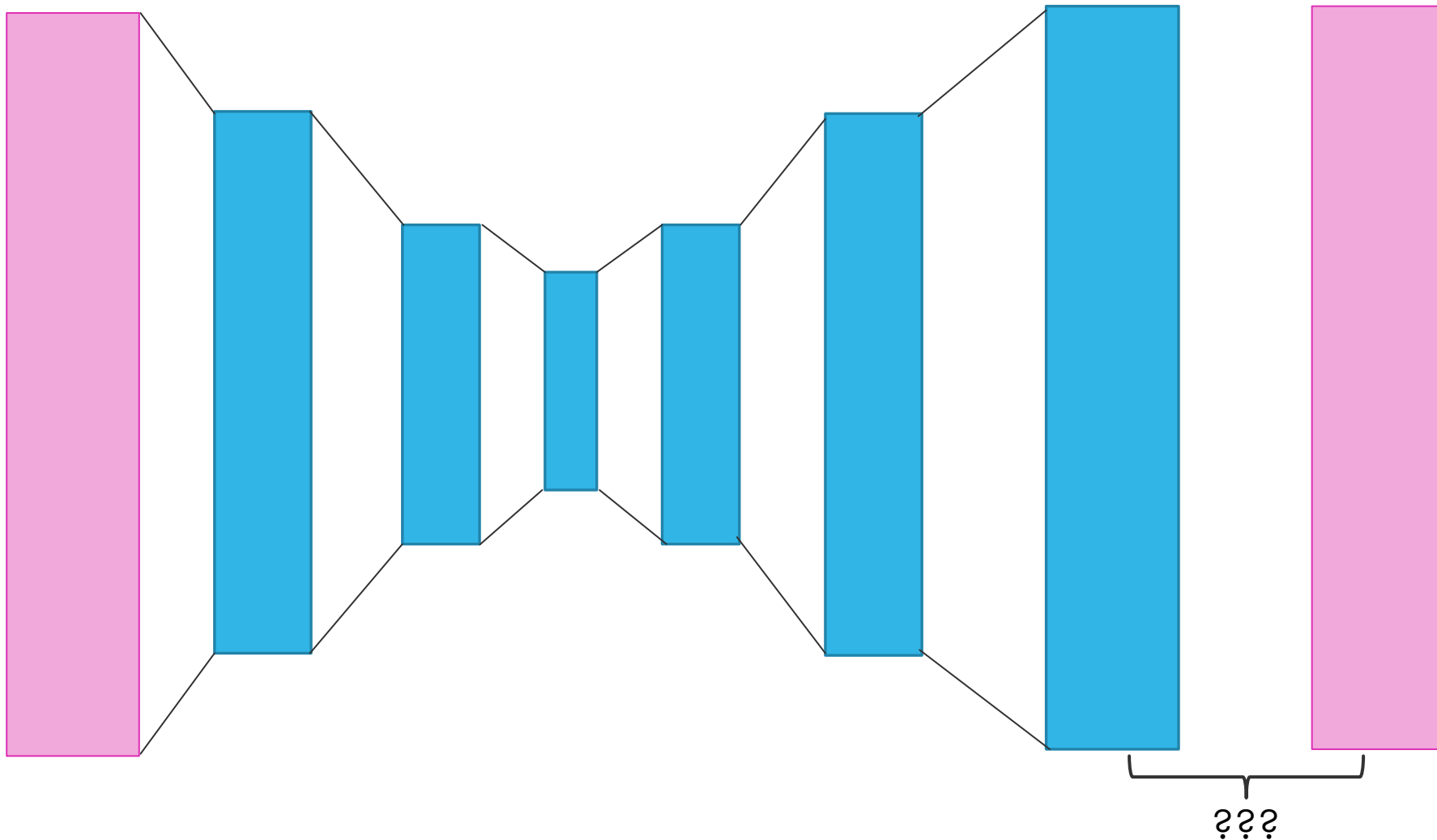




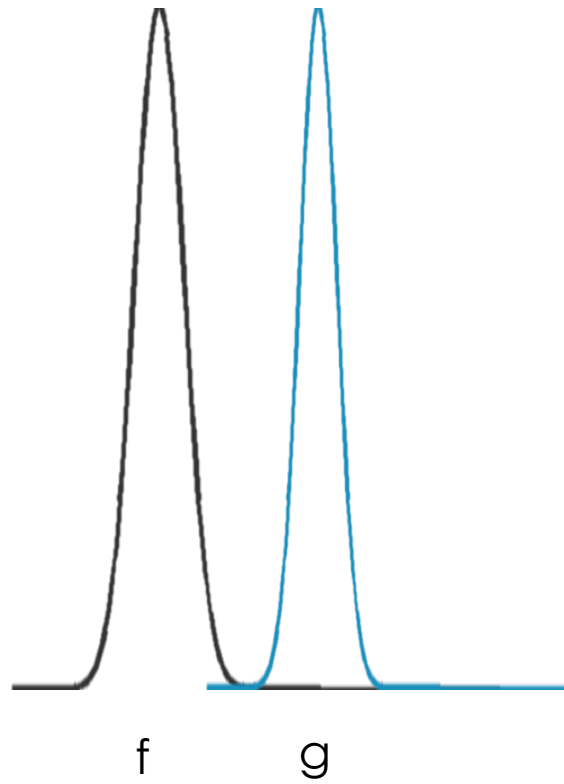
Optimal Transport Principles for Computer Graphics and Machine Learning

Nicolas Bonneel

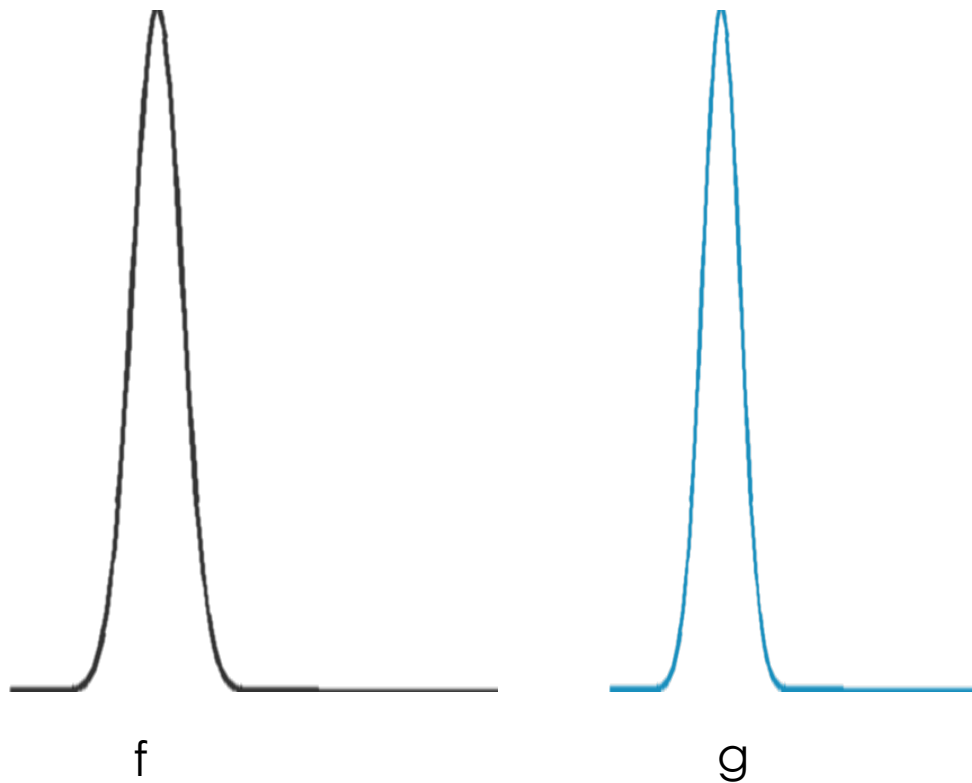
Motivation in neural networks



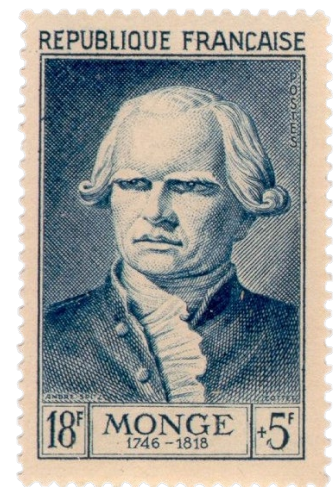
How to compare functions ?



How to compare functions ?

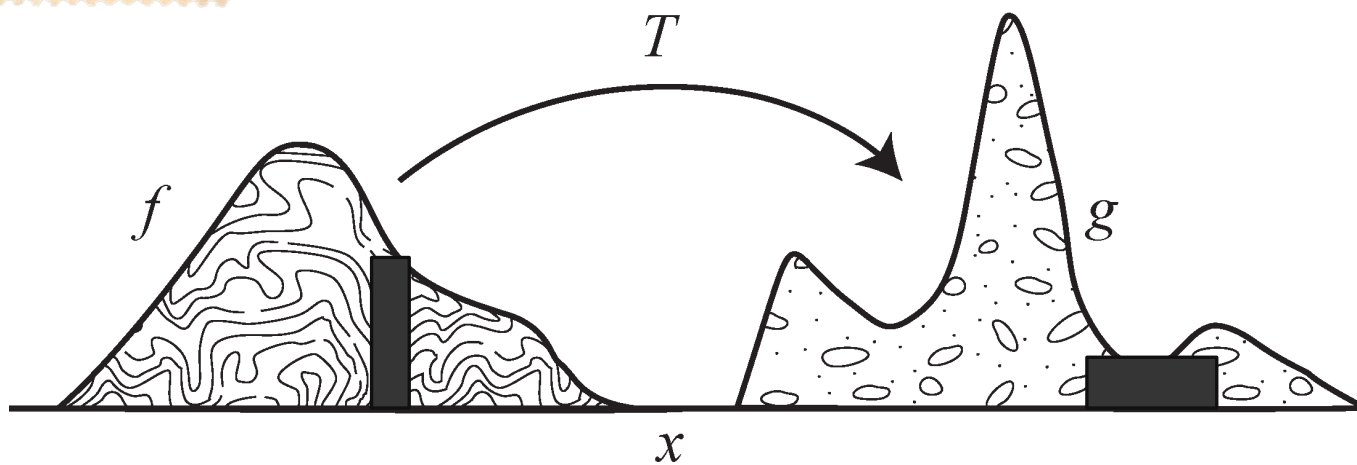


Monge formulation



$$W(f, g) = \min_T \int_X c(x, T(x)) f(x) dx$$

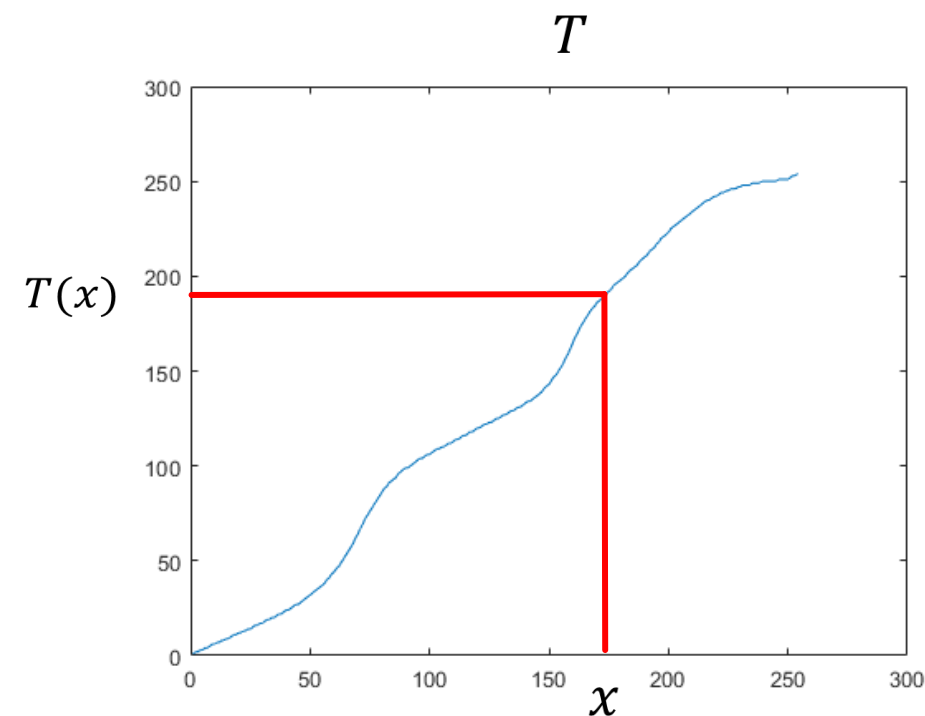
$$\text{s.t. } f(x) = g(T(x)) |\det J_T(x)|$$



Same total mass

“find a good warping between f and g with the change of variable formula”

Monge used $c(x, y) = |x - y|$



Kantorovich formulation



Cost

$$\min_m \sum_i \sum_j c_{i,j} m_{i \rightarrow j}$$

$m_{i \rightarrow j}$ particles will move from i to j

such that:

$$m_{i \rightarrow j} \geq 0$$

Nb of particles is positive !

$$\sum_i m_{i \rightarrow j} = g_j$$

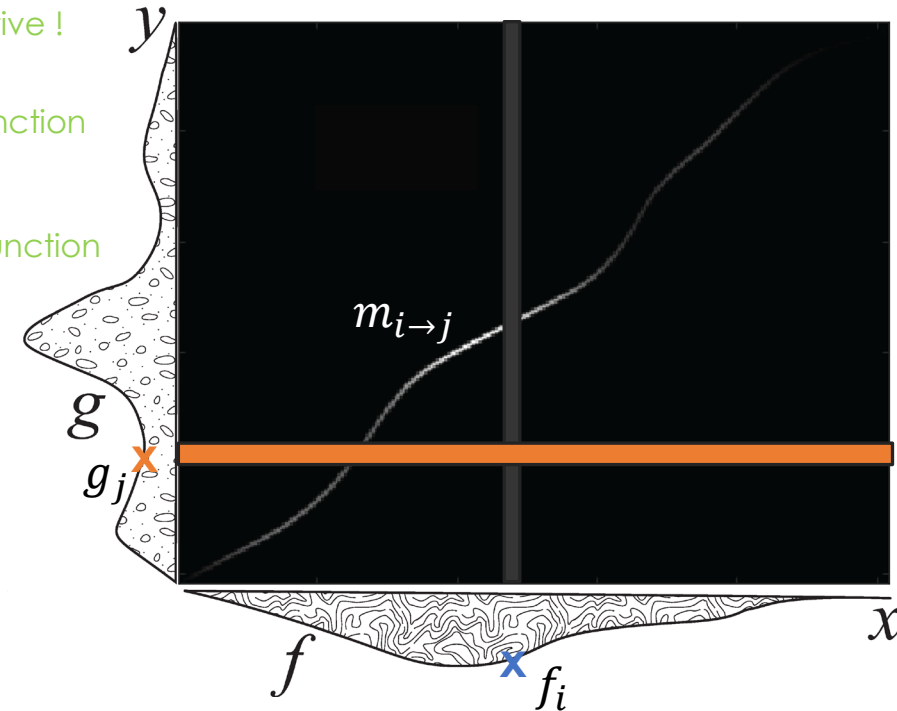
Reconstruct target function

$$\sum_j m_{i \rightarrow j} = f_i$$

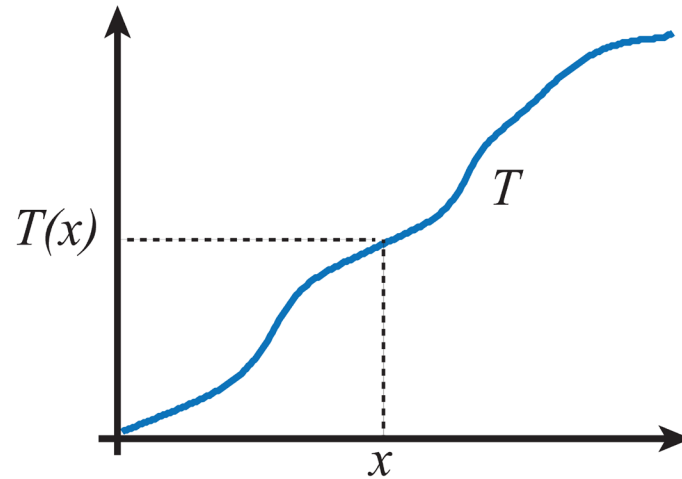
Reconstruct source function

Work for transforming
f into g

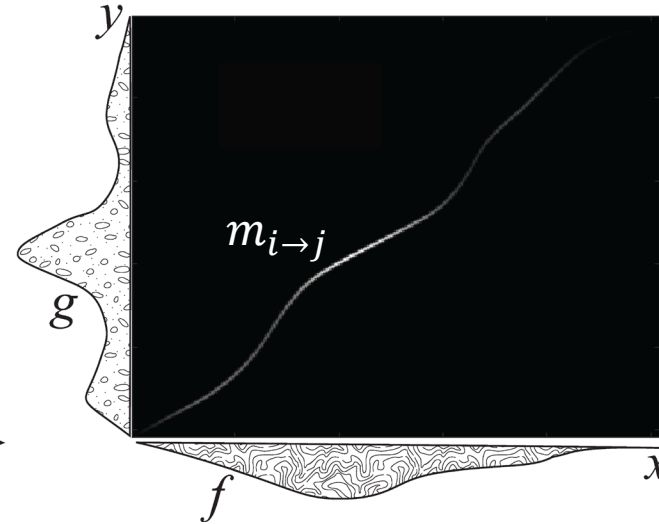
Discretization of the Kantorovich problem
Earth Mover's Distance



Intuition: comparison



- Finds a "transport map"
- Difficult non-linear problem
- May have no solution
(e.g., a Dirac splitting in two)
- Leads to PDEs



- Finds a "transport plan"
- Linear program
- "Always" has solution
(i.e., under reasonable assumptions)
- Also has dual formulation

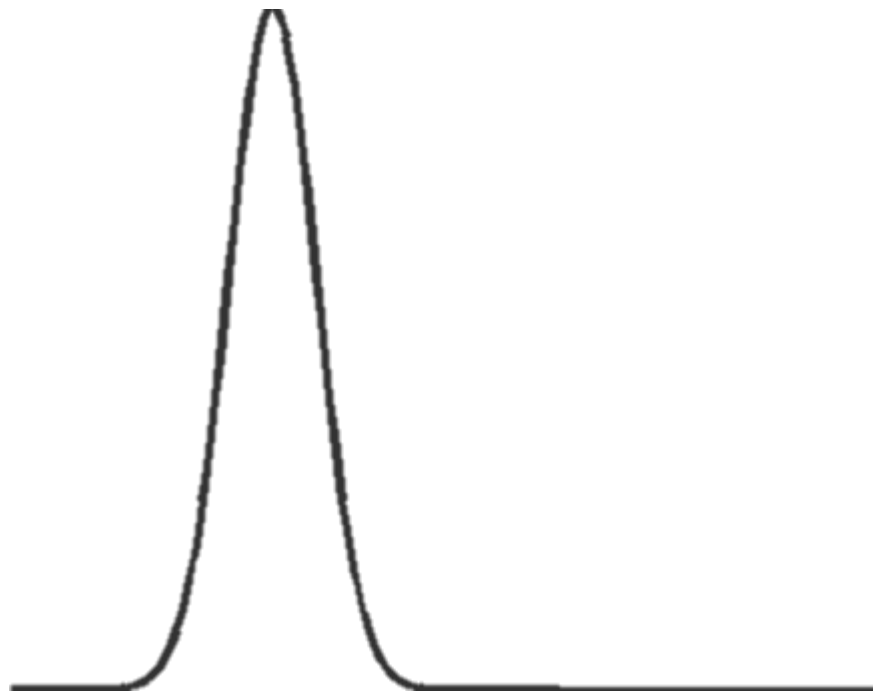
When it exists, the solution is the same.

$$\text{Often, } c(x, y) = \|x - y\|_p^p \Rightarrow W_p^p$$

W_p is a distance



$f(x)$

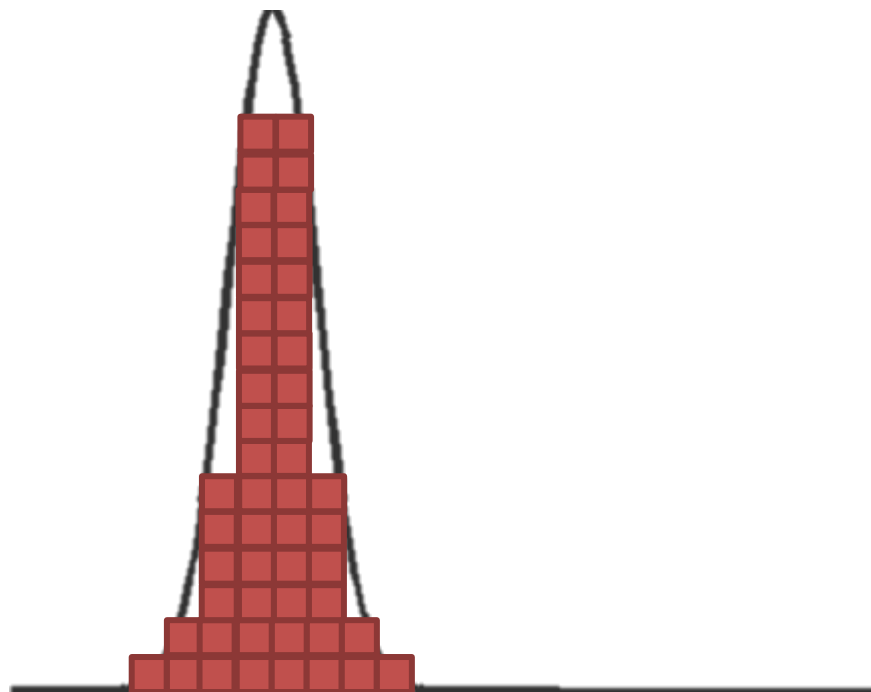


$g(y)$

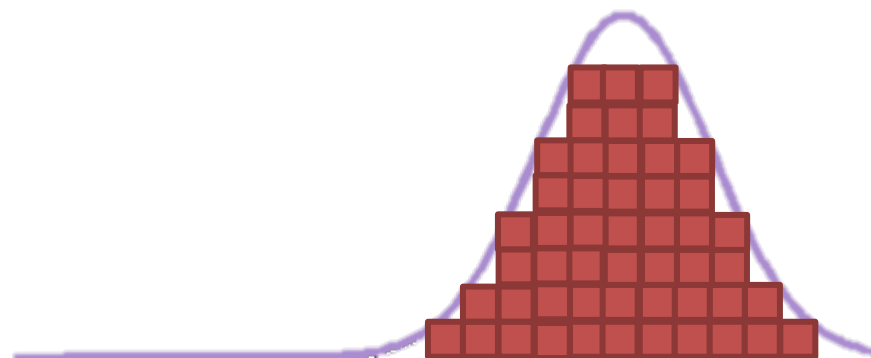


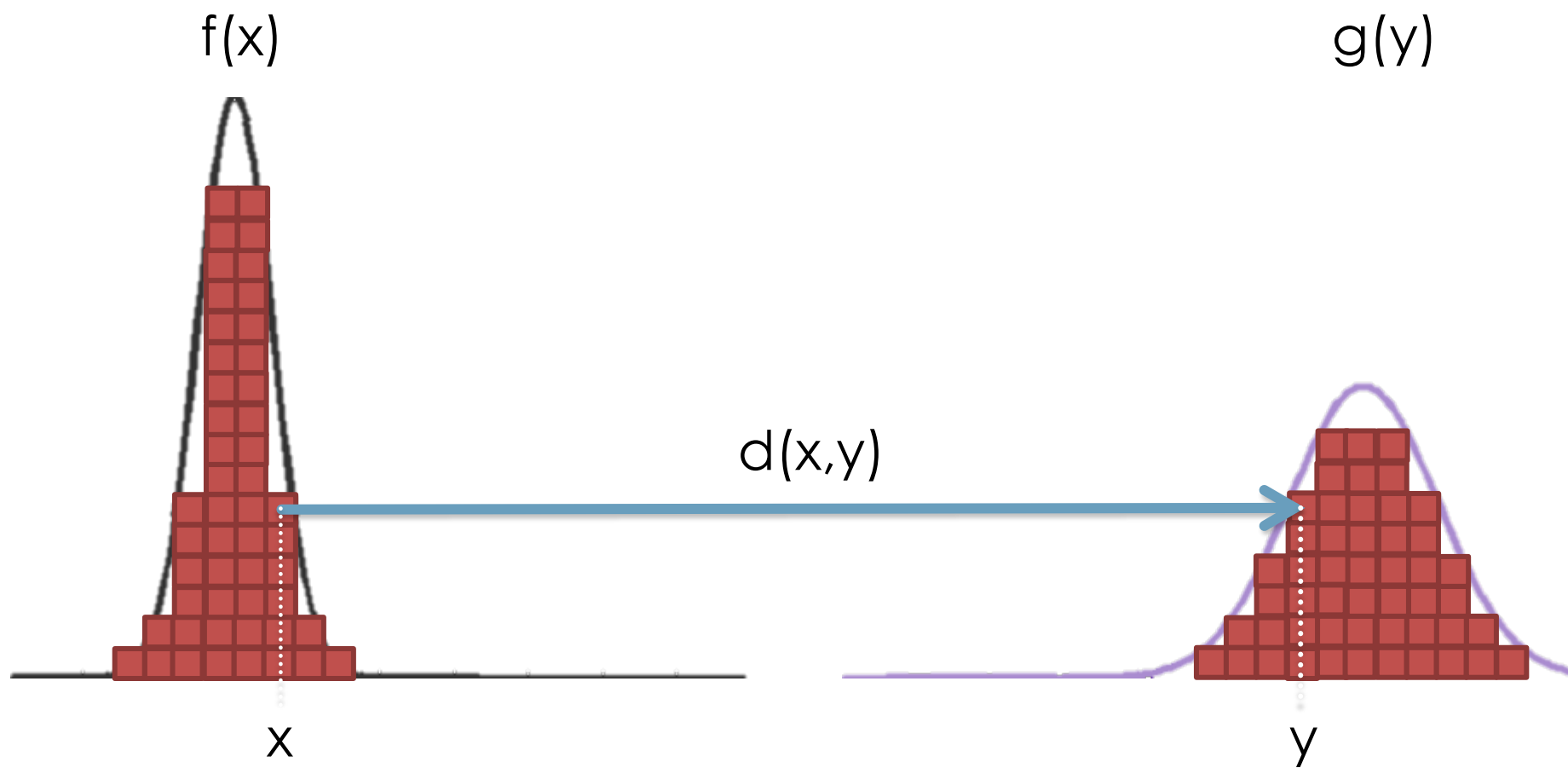


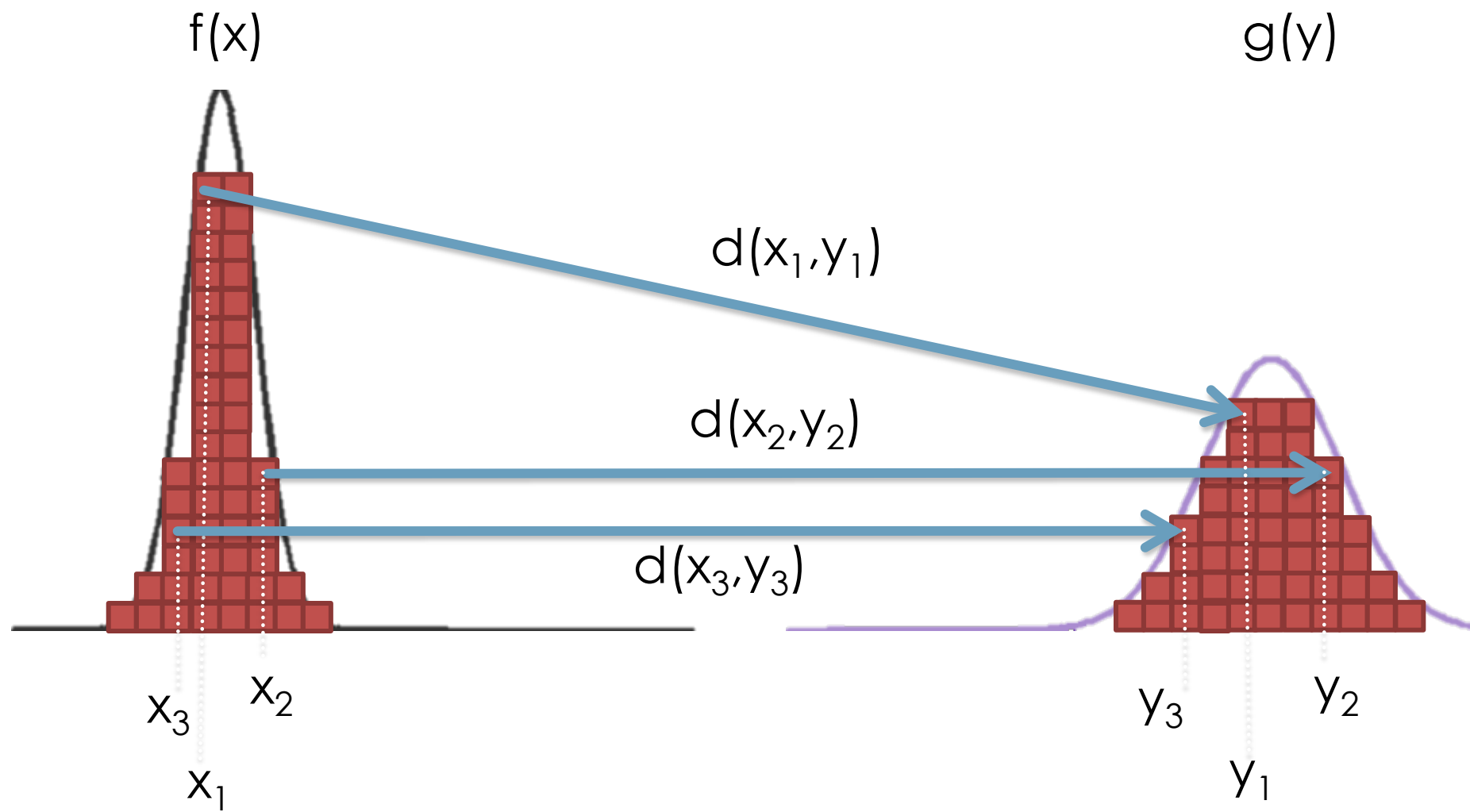
$f(x)$



$g(y)$

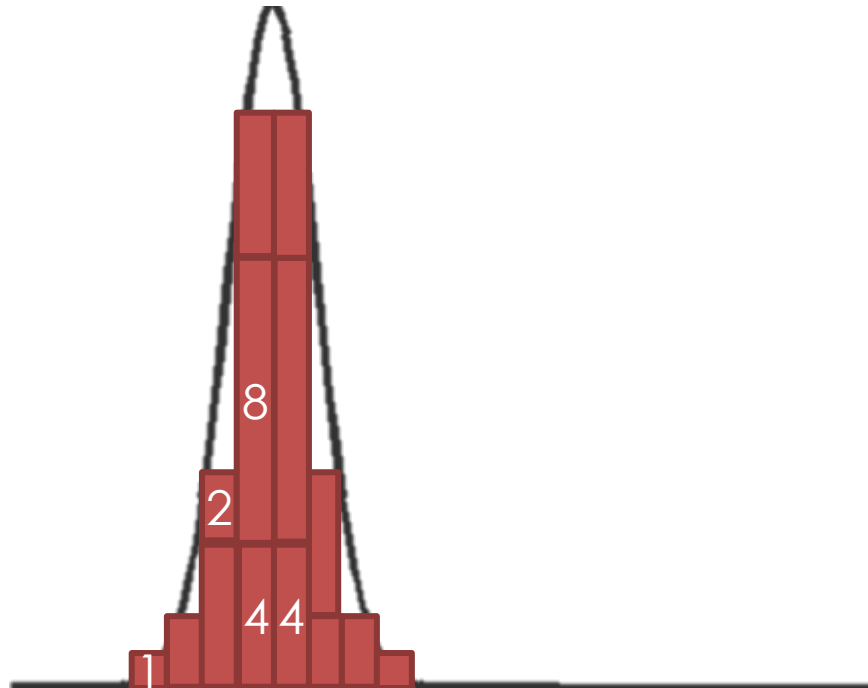




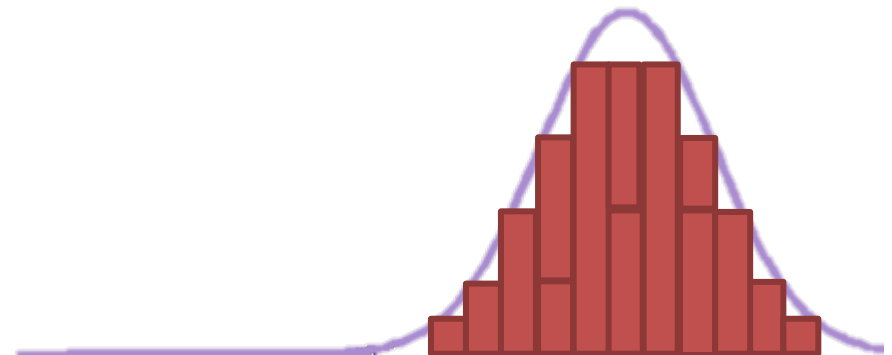




$f(x)$

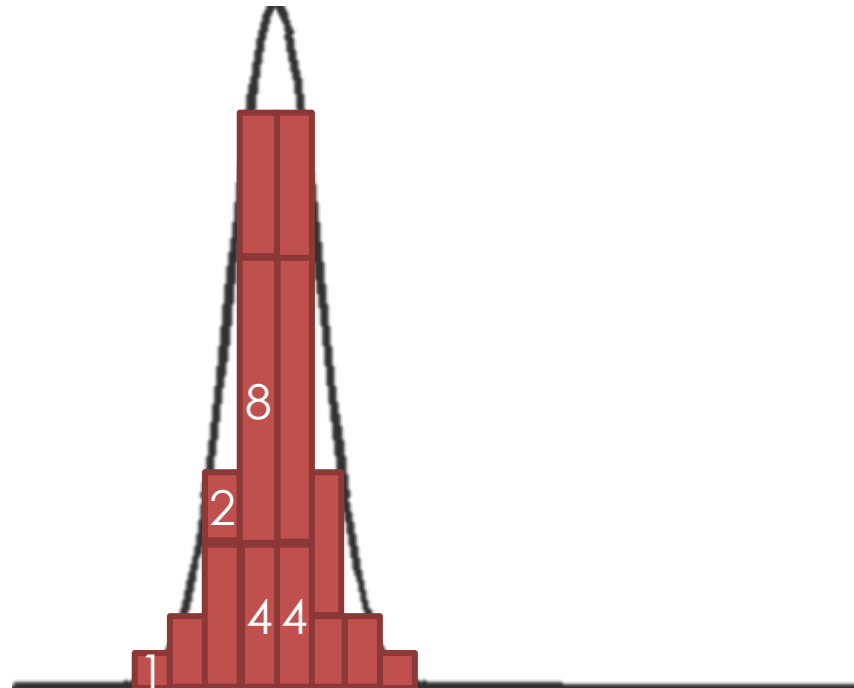


$g(y)$

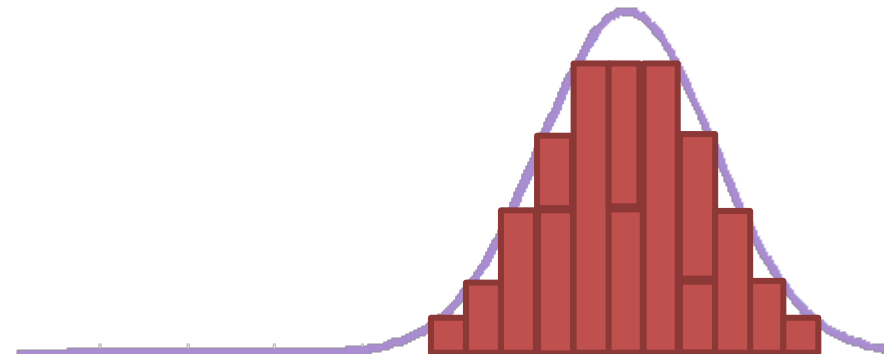




$f(x)$



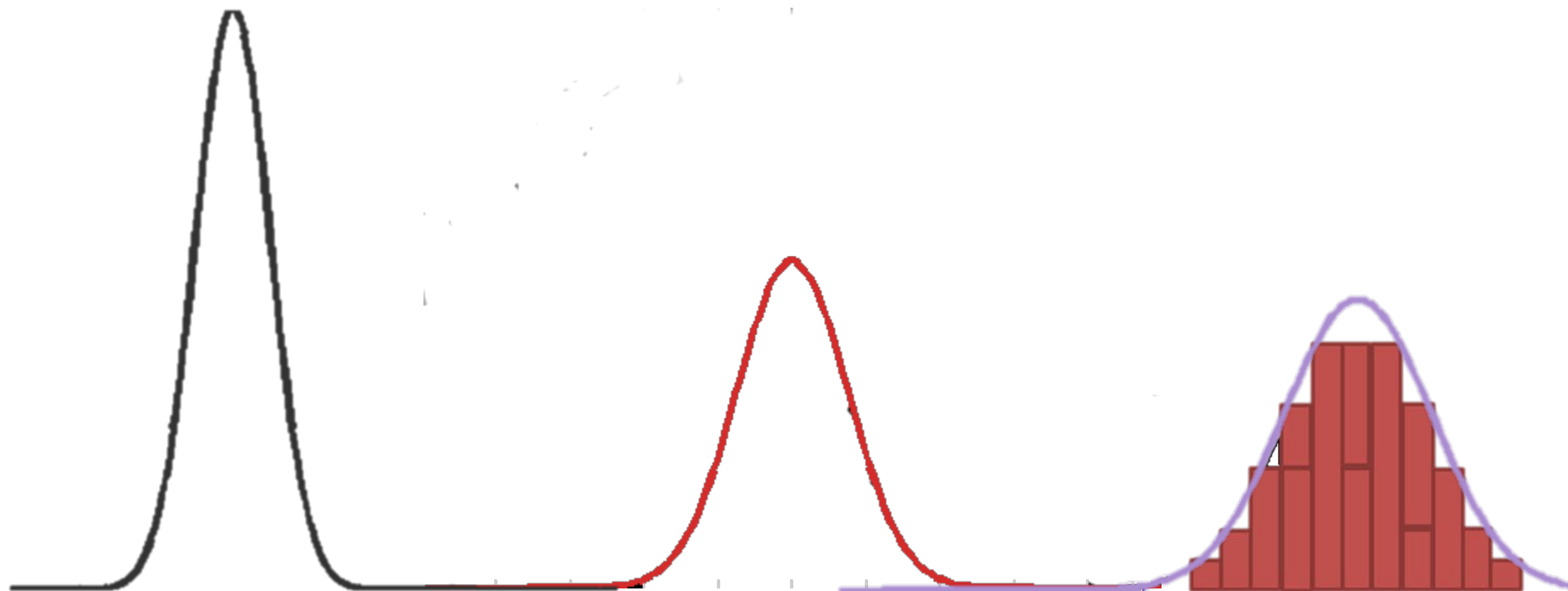
$g(y)$



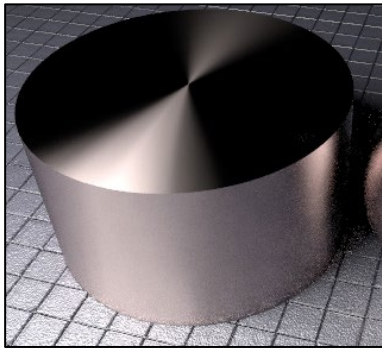


$f(x)$

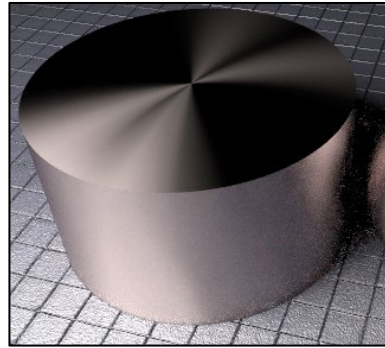
$g(y)$



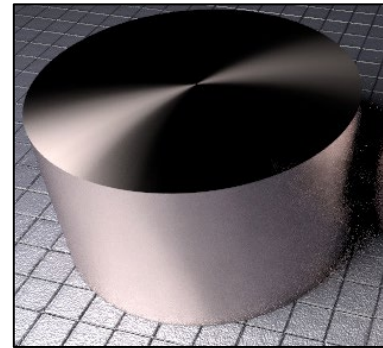
Application: BRDF



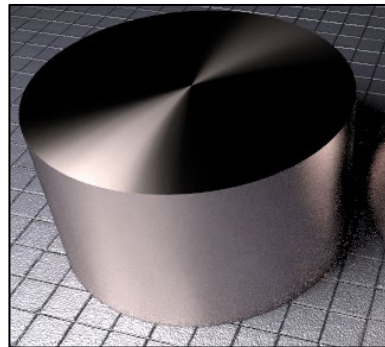
Function A



Linear interpolation

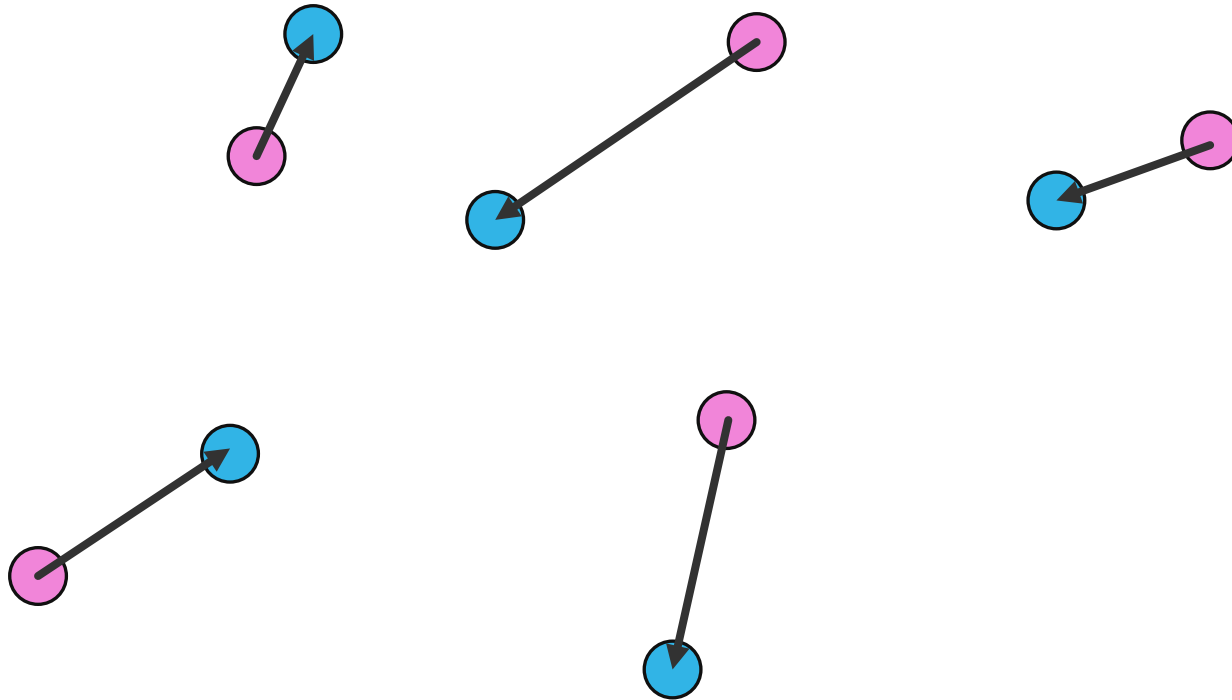


Function B



Displacement interpolation

The assignment problem



Applications to Color Grading



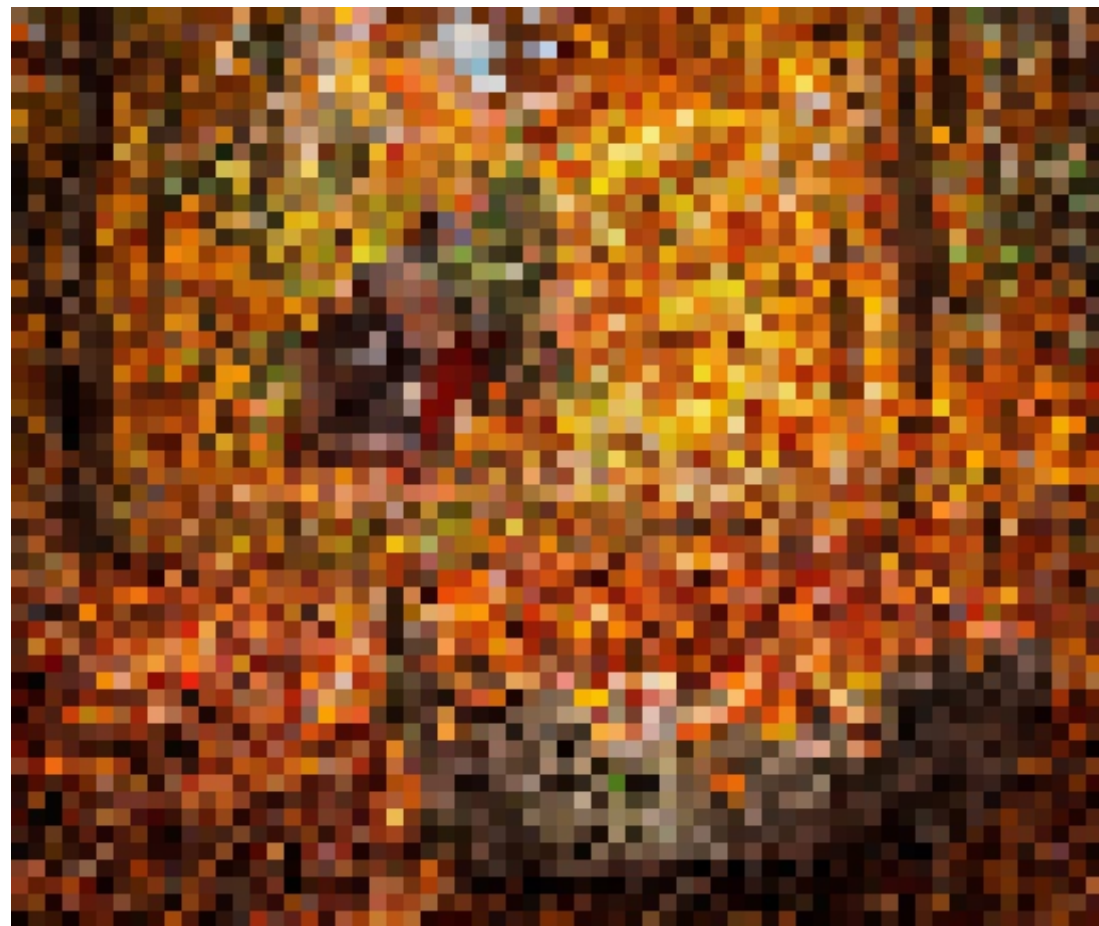
Input photo



Target style

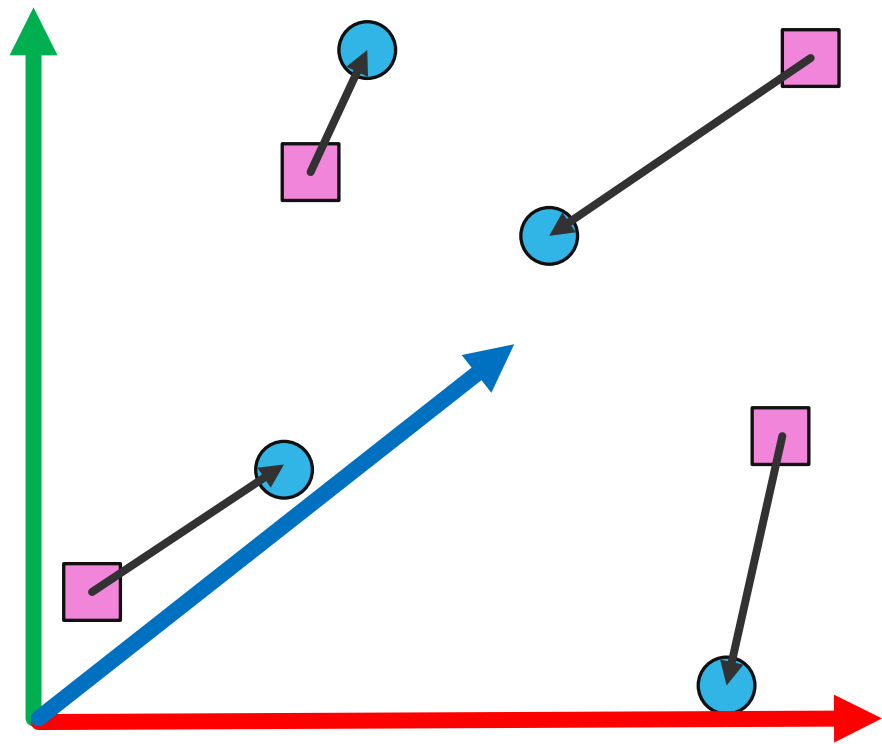


Input photo



Target style

The assignment problem



Pixels

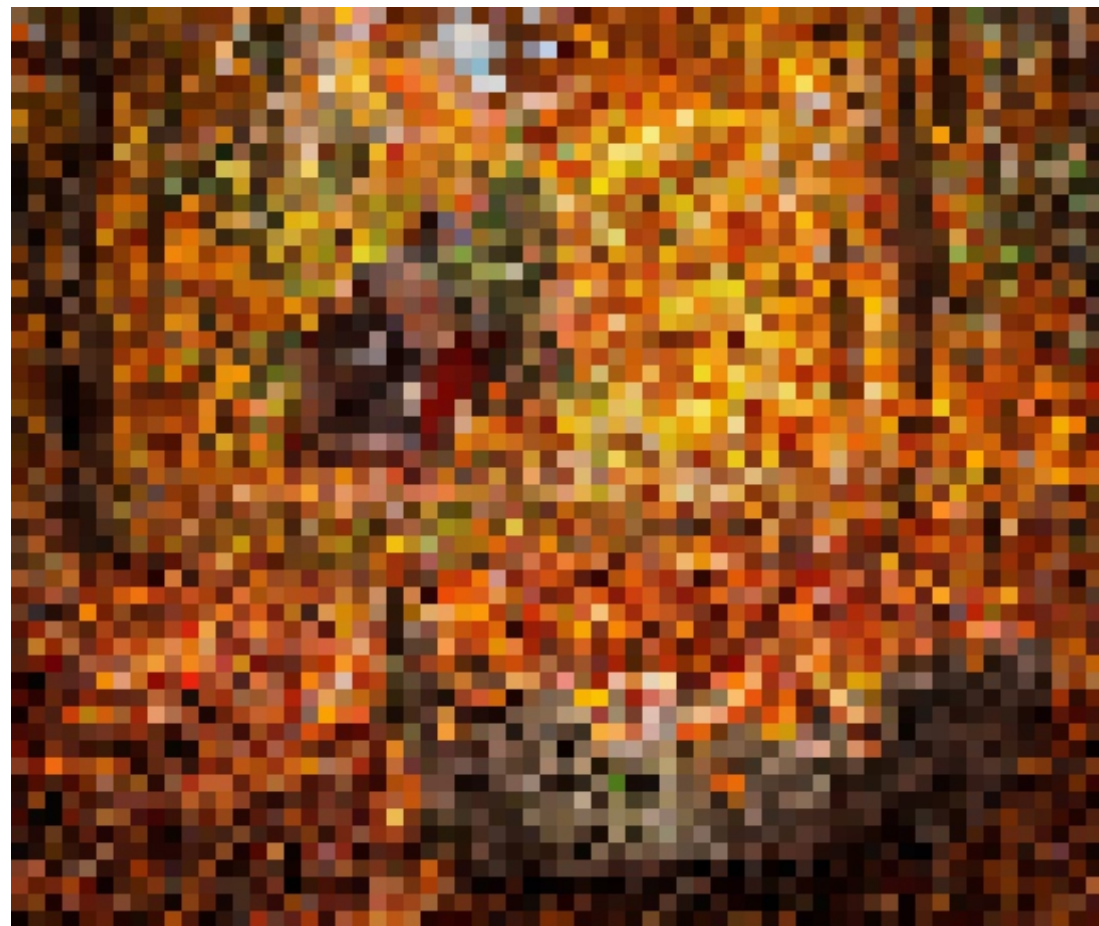


Pixels

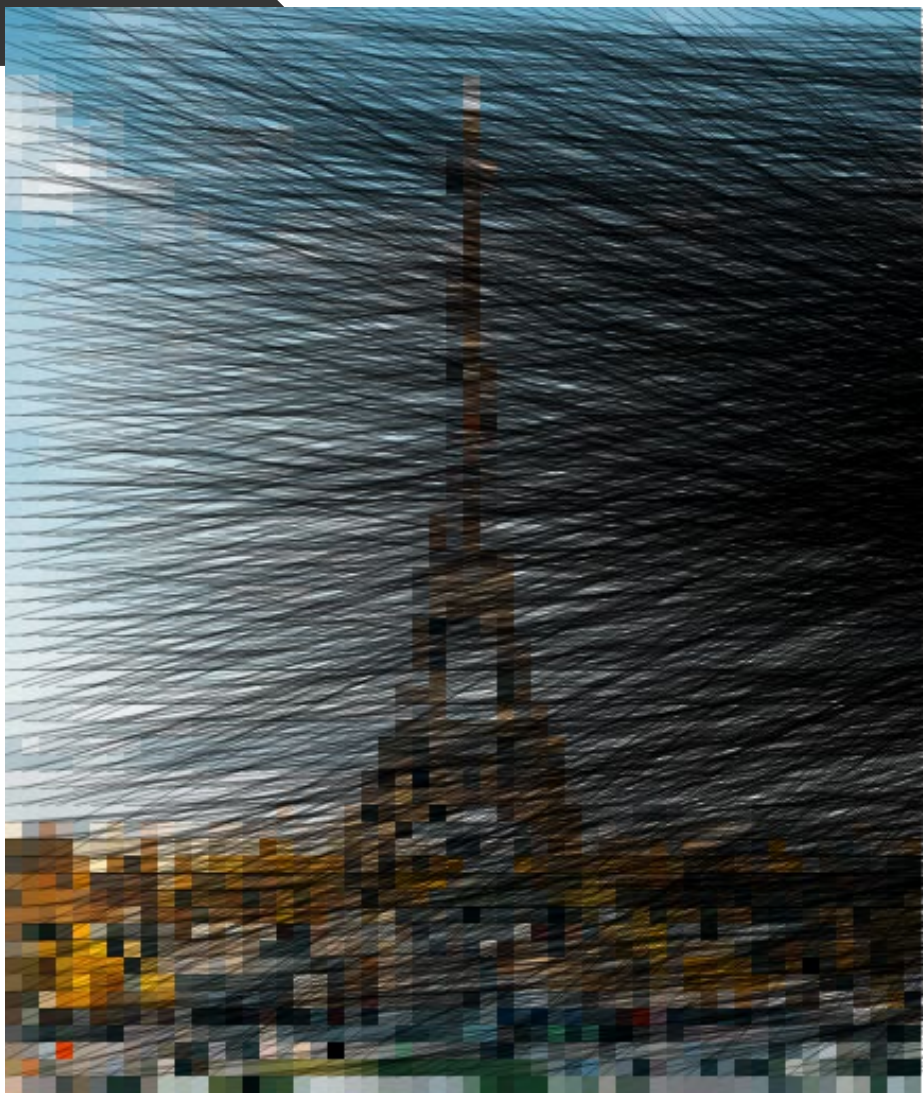




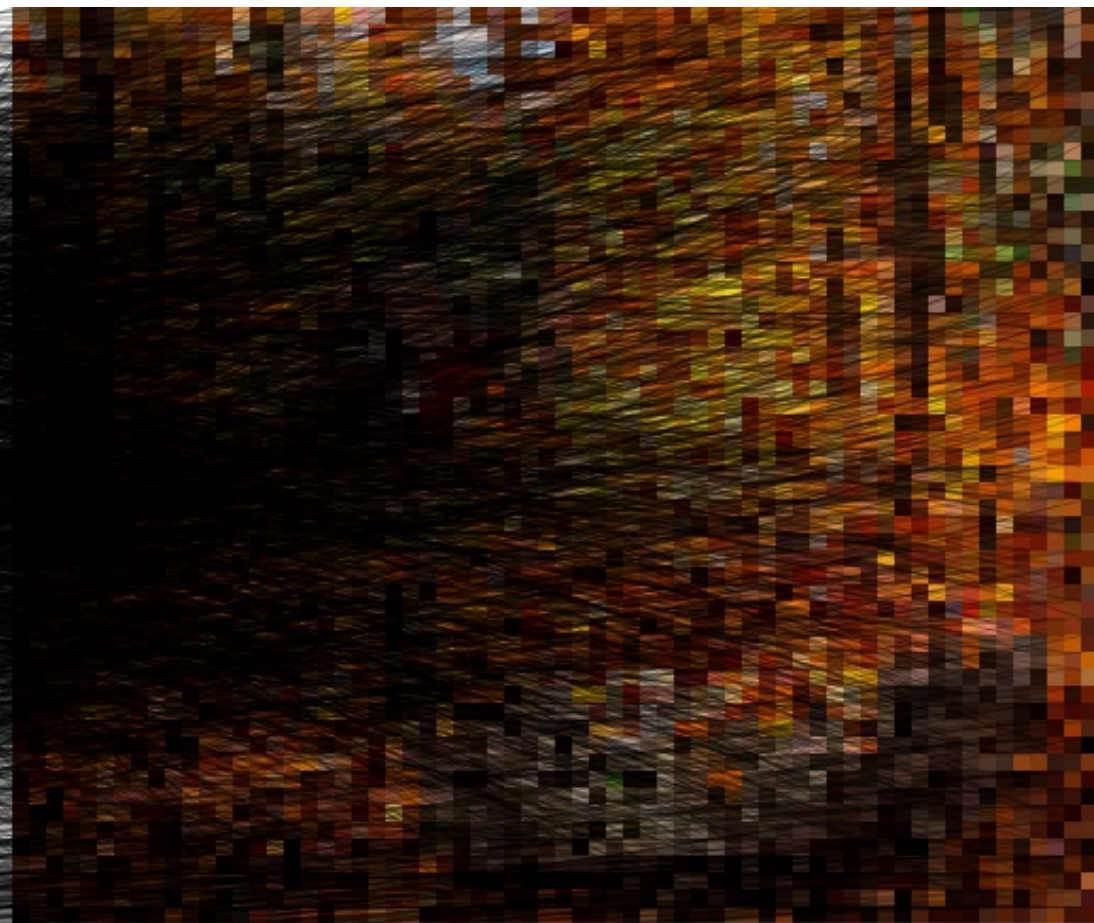
Input photo



Target style



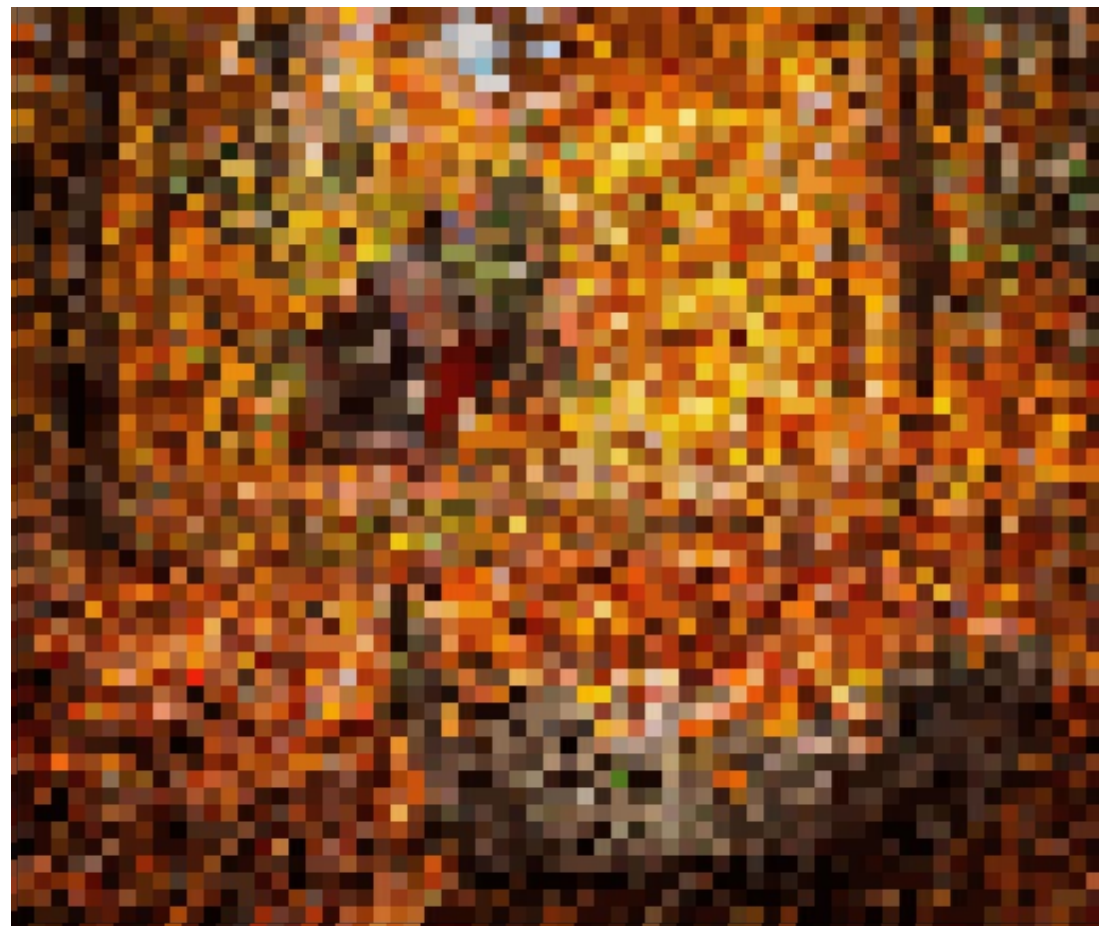
Input photo



Target style



Input photo



Target style





Sliced and Radon Wasserstein Barycenters of Measures

Nicolas Bonneel, Julien Rabin, Gabriel Peyré, Hanspeter Pfister

Journal of Mathematical Imaging and Vision (2014)



Simple cases

- Transport 1 Gaussian \leftrightarrow 1 Gaussian
- Transport = translation + scaling
- Transport 1D function \leftrightarrow 1D function

Optimal transport is simple for Gaussians

- Optimal transport and barycenters trivially solved for
 - Gaussian distributions with $c(x, y) = \|x - y\|^2$
 - $W_2^2(\mathcal{N}_0, \mathcal{N}_1) = \text{tr}(\Sigma_0 + \Sigma_1 - 2\Sigma_{0,1}) + \|\mu_0 - \mu_1\|^2$ with $\Sigma_{0,1} = \left(\Sigma_0^{-\frac{1}{2}}\Sigma_1\Sigma_0^{-\frac{1}{2}}\right)^{1/2}$
 - $T(x) = \Sigma_{0,1}x$
 - Barycenter: $\mathcal{N}(\mu, \Sigma)$ with $\mu = \sum_k \lambda_k \mu_k$ and iterations

$$\Sigma^{(n+1)} = \sum_k \lambda_k \left(\sqrt{\Sigma^{(n+1)}} \Sigma_k \sqrt{\Sigma^{(n+1)}} \right)^{1/2}$$

Optimal transport is simple in 1D

- Continuous case with density, convex cost, $\mu = f dx$, $\nu = g dy$

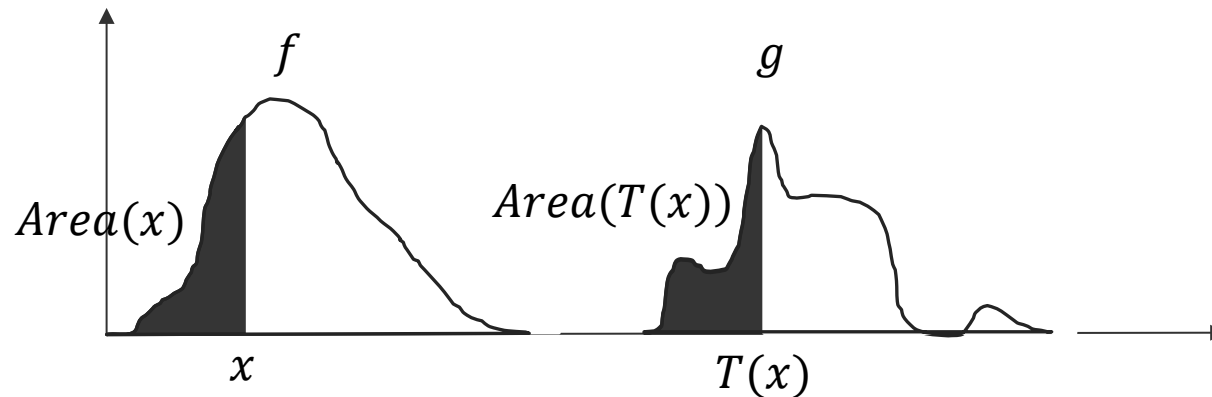
- Need: $\int_{-\infty}^x f(x)dx = \int_{-\infty}^{T(x)} g(x)dx$

$$T = G^{-1} \circ F$$

with $F(x) = \int_{-\infty}^x f(x)dx$ and $G(x) = \int_{-\infty}^x g(x)dx$

Generalize G^{-1} : $G^{-1}(y) = \min_x \{y = G(x)\}$ \longrightarrow

Quantile function: e.g.:
“what salary corresponds to the first percentile”





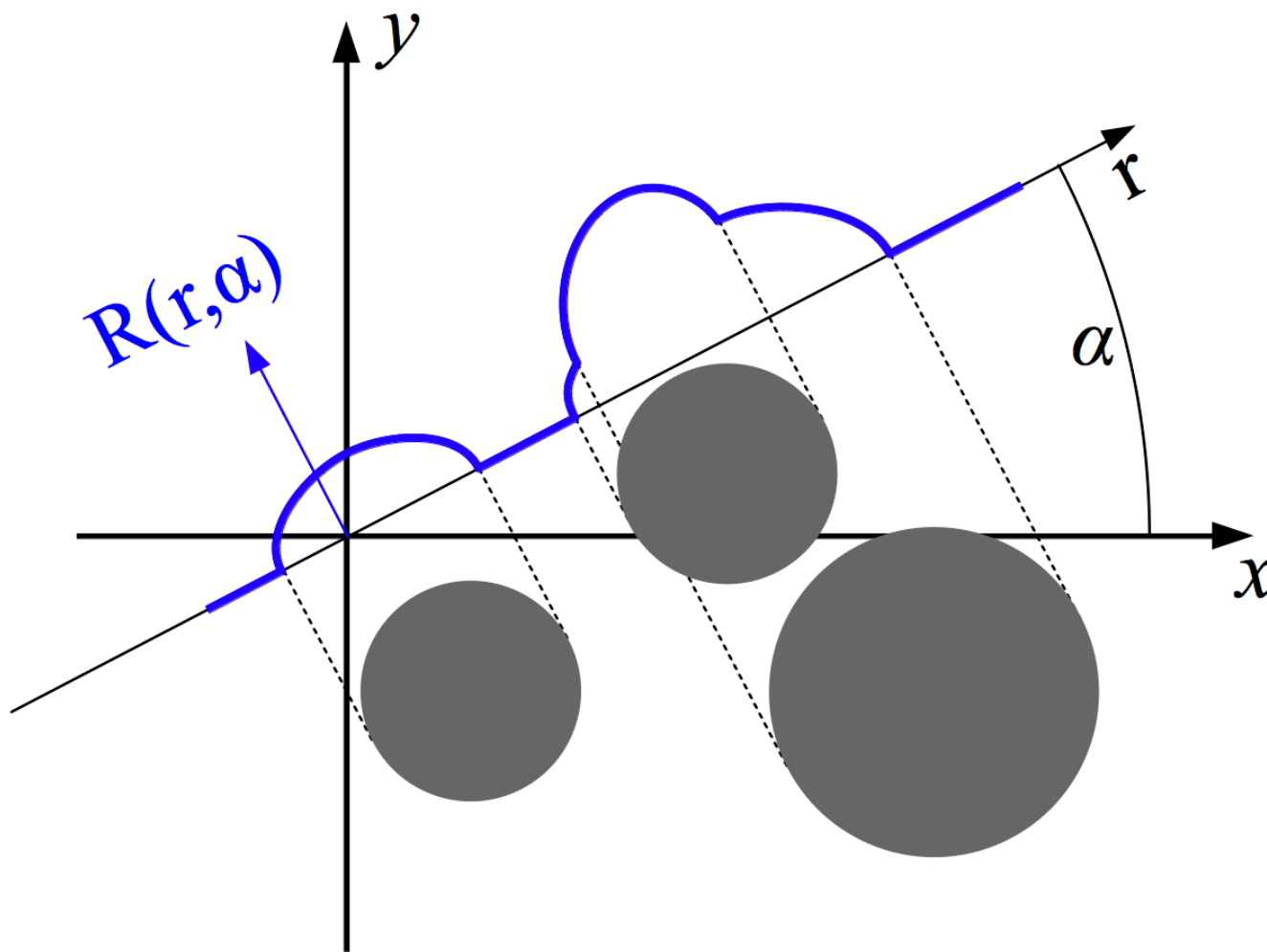
1D Case

OT Map: $T = G^{-1} \circ F$

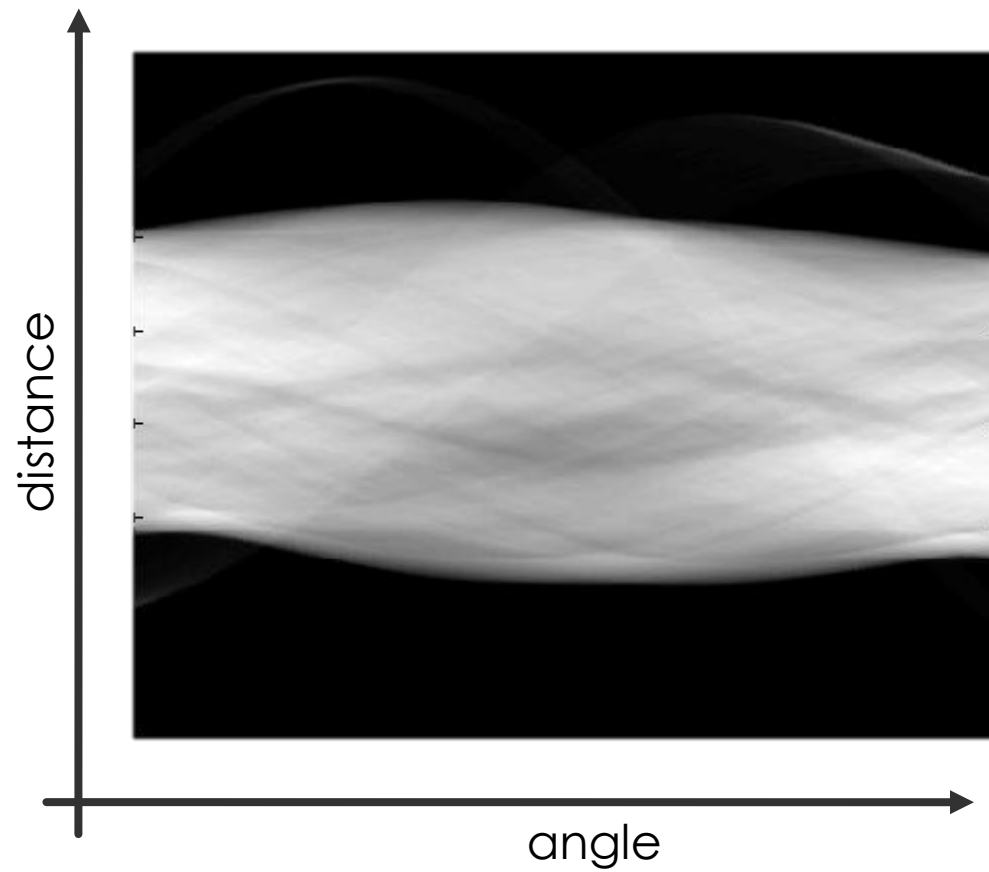
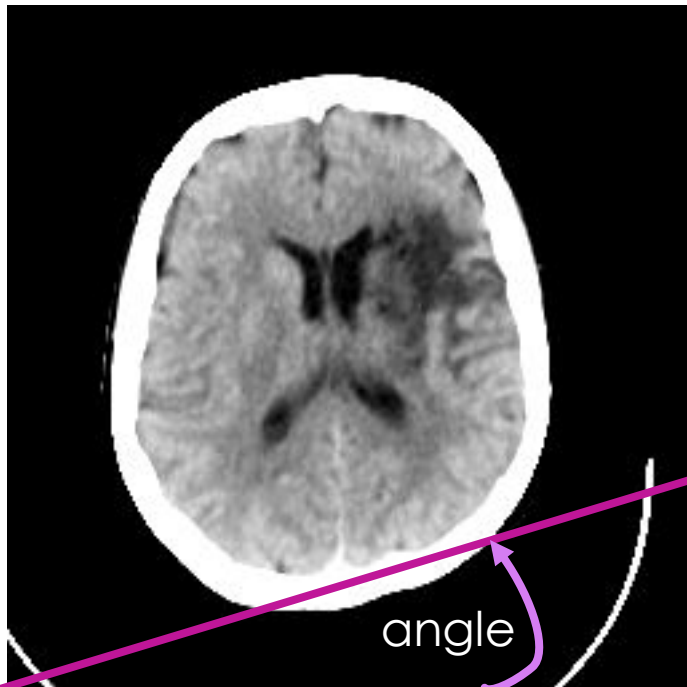
OT cost: $\int_0^1 c(F^{-1}(t) - G^{-1}(t)) dt$

Interpolation: $F_{interp}^{-1}(x) = \sum_i \alpha_i F_i^{-1}(x)$

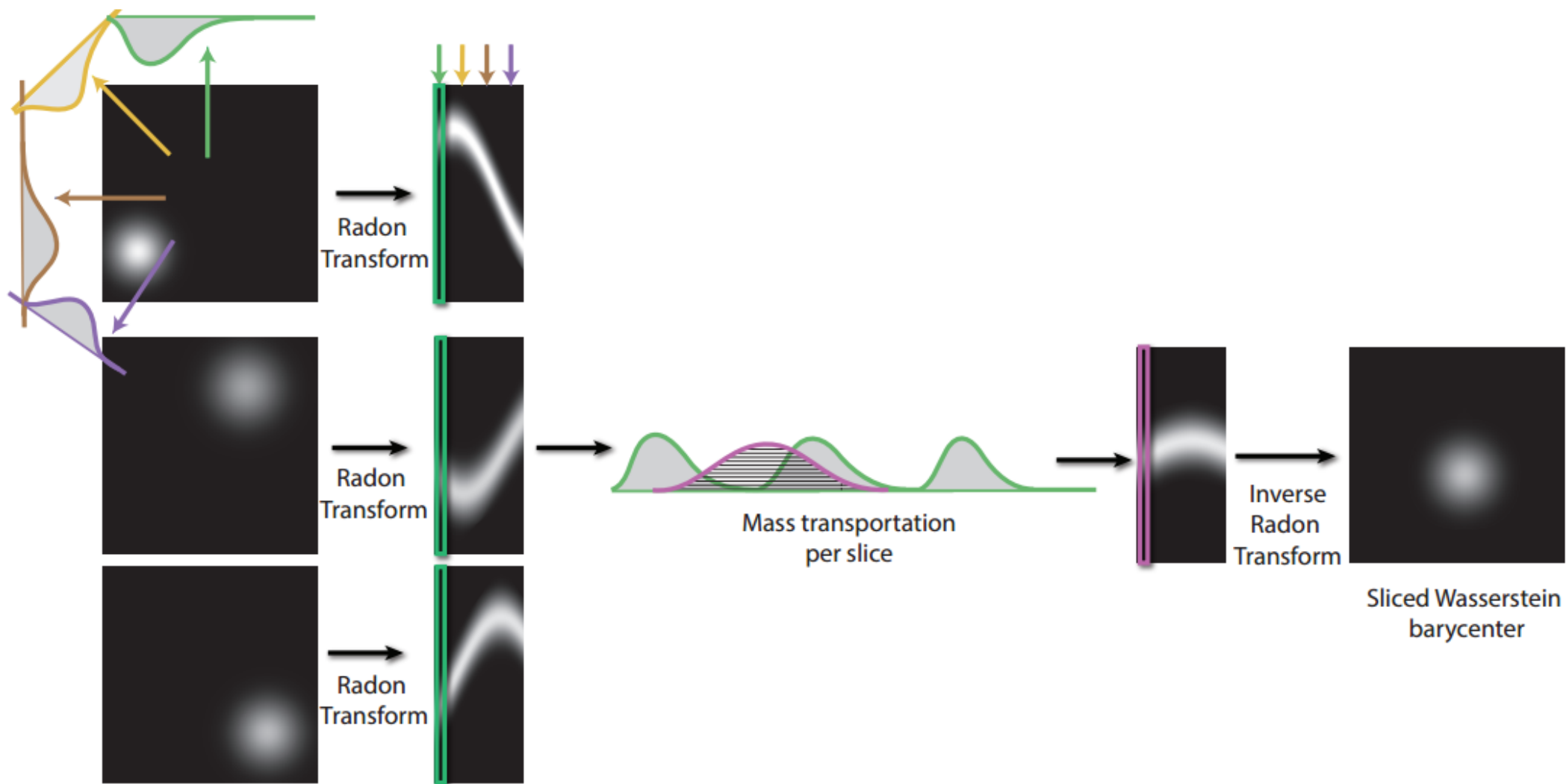
Radon transform

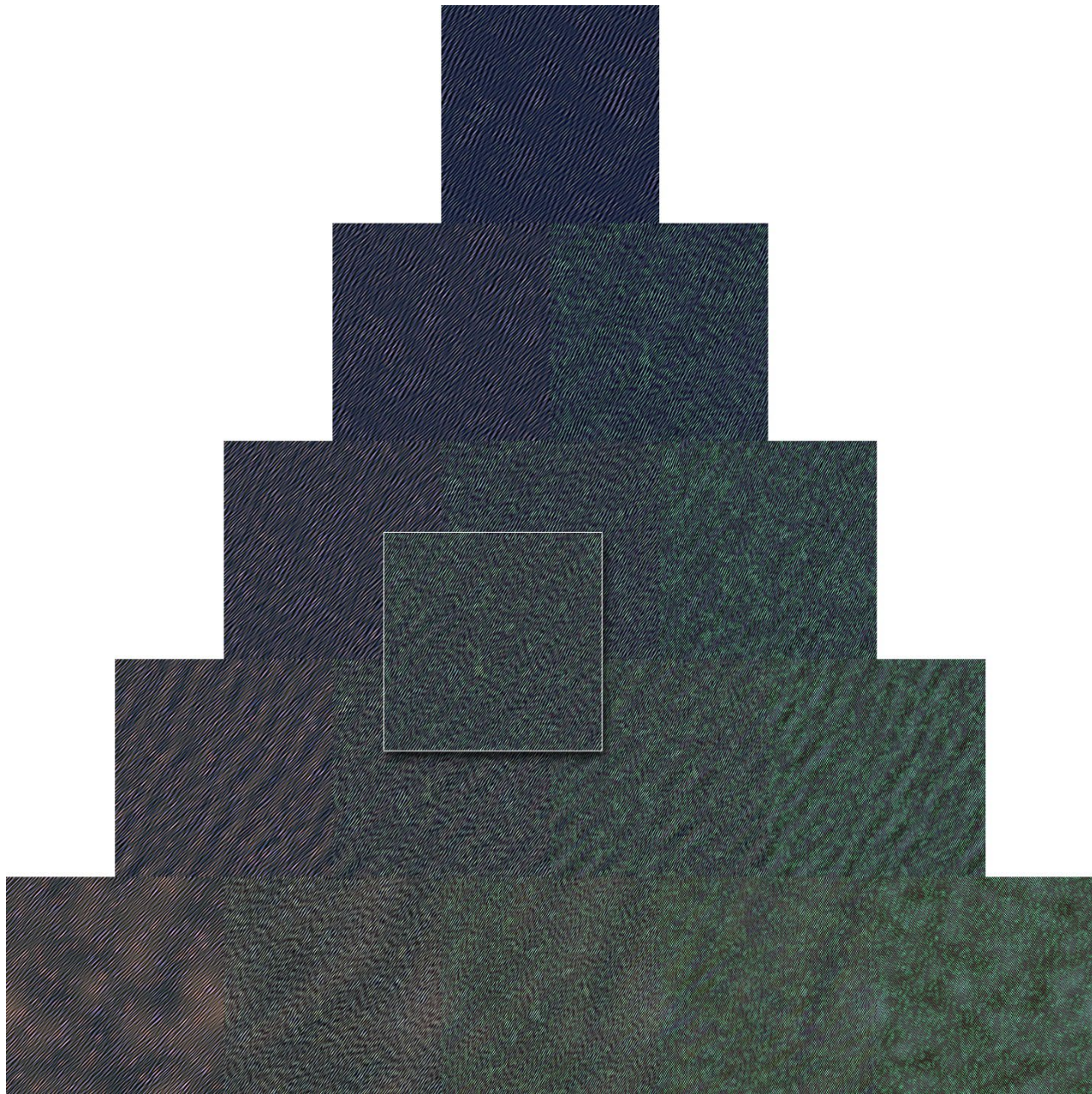
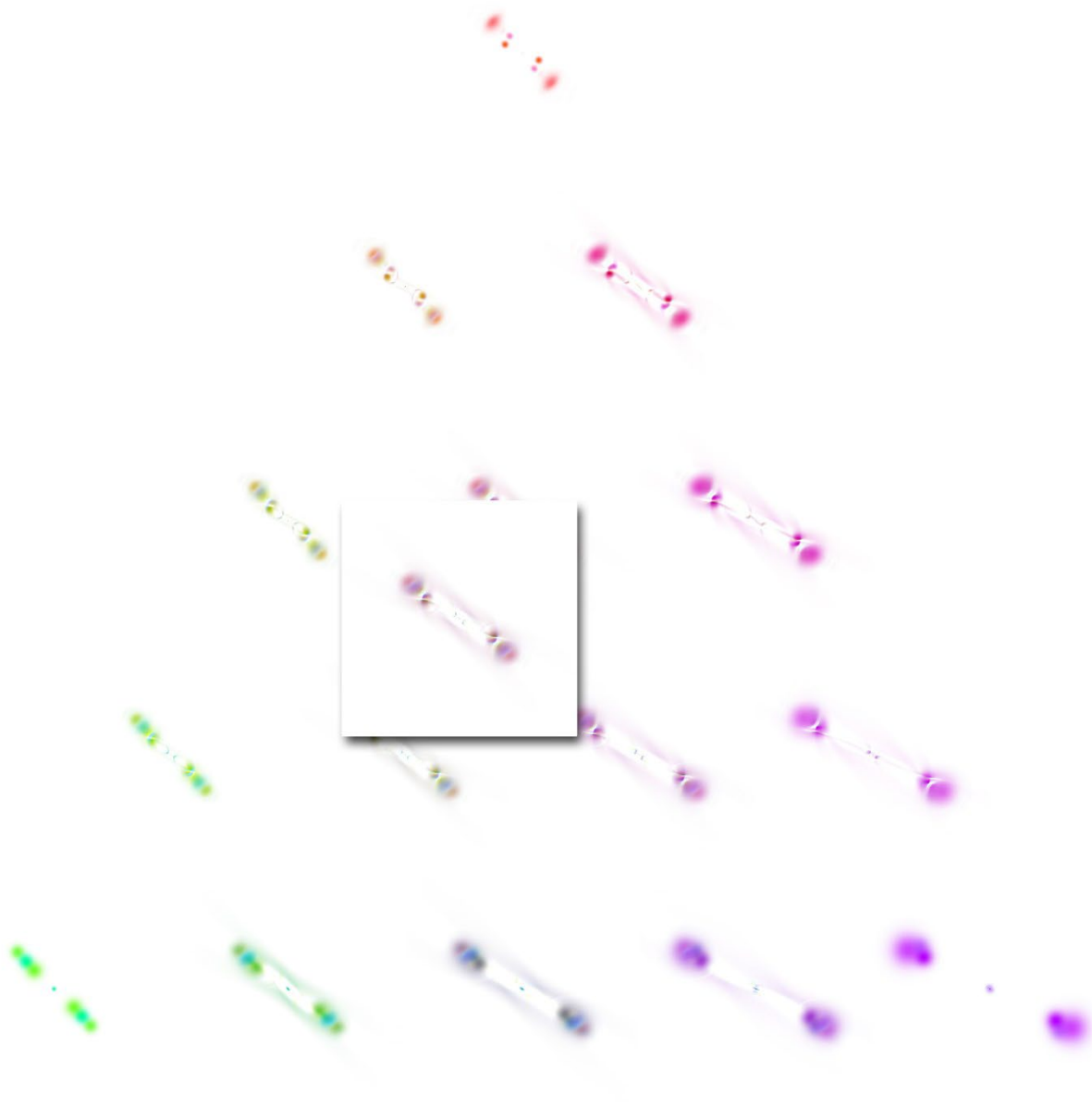


Radon transform



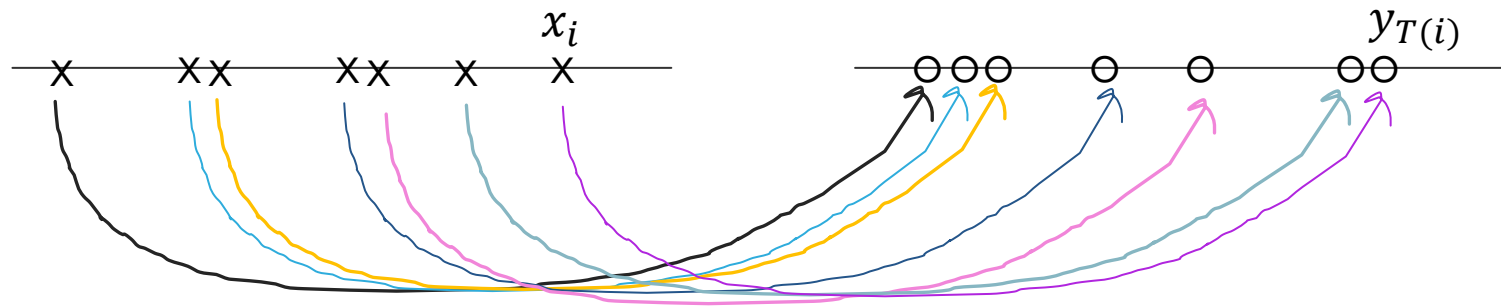
Method





1D Case, discrete

- Discrete case, $\mu = \sum_{i=1}^n \delta_{x_i}$, $\nu = \sum_{i=1}^n \delta_{y_i}$ (same for interpolating between more than 2 measures)
- Optimal transport for convex cost = pairing sorted samples





Sliced Wasserstein Distance

► For discrete high-dimensional distributions $\mu = \sum_{i=1}^n \delta_{x_i}$ and $\nu = \sum_{i=1}^n \delta_{y_i}$

Consider energy

$$SW(\mu, \nu) = \int_S W_2^2(\text{proj}(\mu, \omega), \text{proj}(\nu, \omega)) d\omega$$

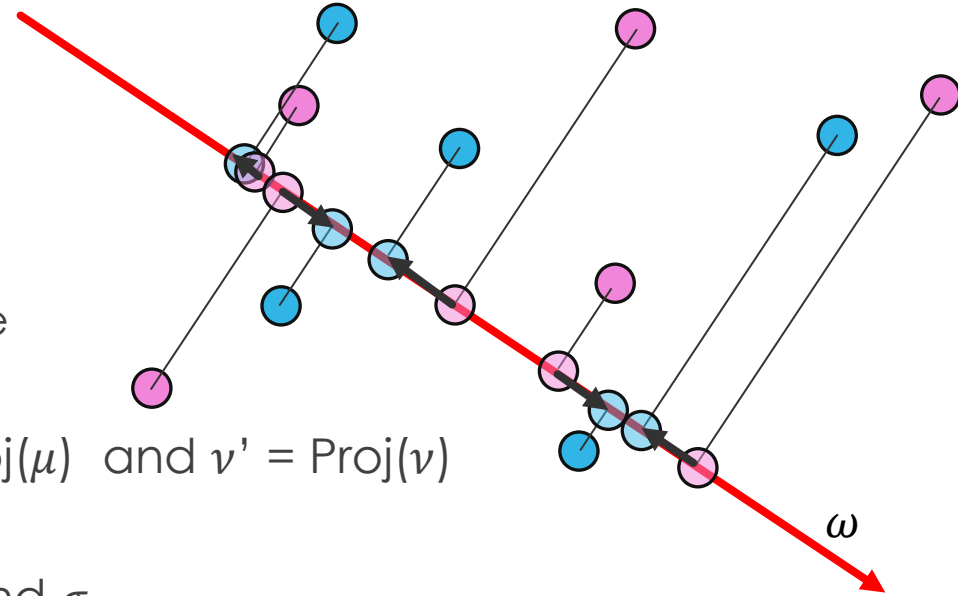
Where $\text{proj}(\mu, \omega)$ is the 1-d distribution : $\text{proj}(\mu, \omega) = \sum_i \delta_{\langle x_i, \omega \rangle}$ (same for ν)

And W_2^2 computes the 1-d squared Wasserstein distance

Sliced Wasserstein Distance

- Take a uniform random direction ω
 - $\omega \leftarrow (\mathcal{N}(0,1), \mathcal{N}(0,1), \mathcal{N}(0,1))$ and normalize
- Project samples of μ and ν on ω : $\mu' = \text{Proj}(\mu)$ and $\nu' = \text{Proj}(\nu)$
- Sort μ' and ν' , i.e, find permutations σ_μ and σ_ν
- To compute the Sliced Wasserstein Distance:

$$d^2 \leftarrow d^2 + \sum_i \left| \langle x_{\sigma_\mu(i)}, \omega \rangle - \langle y_{\sigma_\nu(i)}, \omega \rangle \right|^2$$



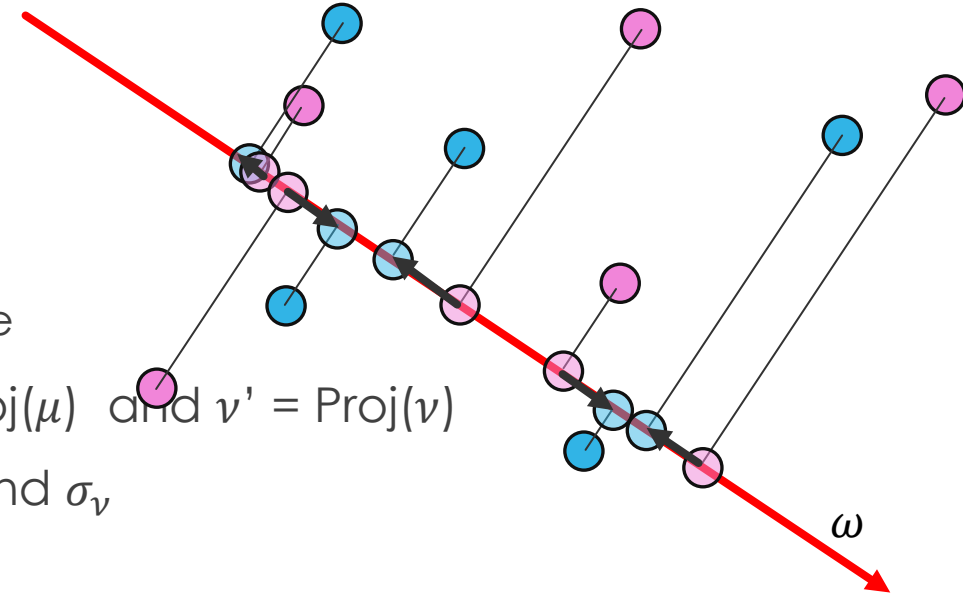


Gradient flow

- Given an energy (e.g., the Sliced Wasserstein Distance), you may try to minimize it with gradient descent
- Each step of such a gradient descent is a “gradient flow”
- This “moves” the input distribution towards the target (the target is such that the energy is 0)

Sliced Wasserstein Gradient Flow

- Take a uniform random direction ω
 - $\omega \leftarrow (\mathcal{N}(0,1), \mathcal{N}(0,1), \mathcal{N}(0,1))$ and normalize
- Project samples of μ and ν on ω : $\mu' = \text{Proj}(\mu)$ and $\nu' = \text{Proj}(\nu)$
- Sort μ' and ν' , i.e, find permutations σ_μ and σ_ν
- Update μ by $x_{\sigma_\mu(i)} \leftarrow x_{\sigma_\mu(i)} + \left(\langle y_{\sigma_\nu(i)}, \omega \rangle - \langle x_{\sigma_\mu(i)}, \omega \rangle \right) \omega$
 - Corresponds to moving particles in the direction ω



Sliced Wasserstein Distance



f

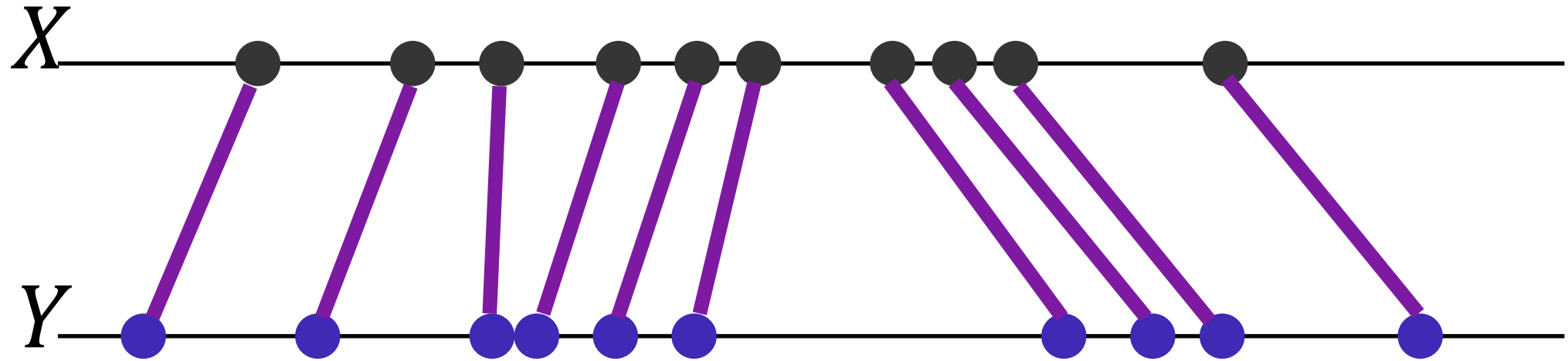


g





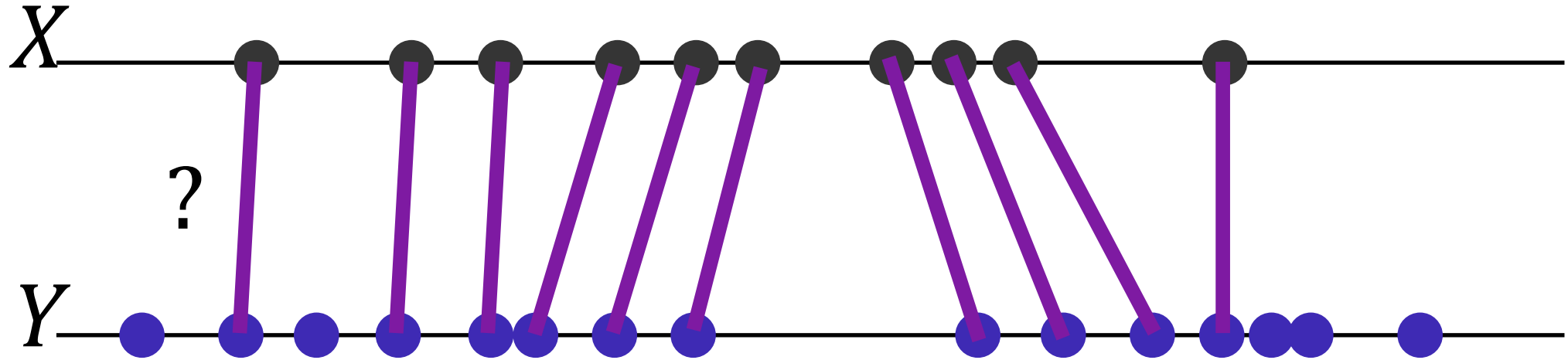
1-d Linear Assignment Problem is trivial*



*assuming the cost c is a convex function of $|x-y|$

Partial optimal assignment ?

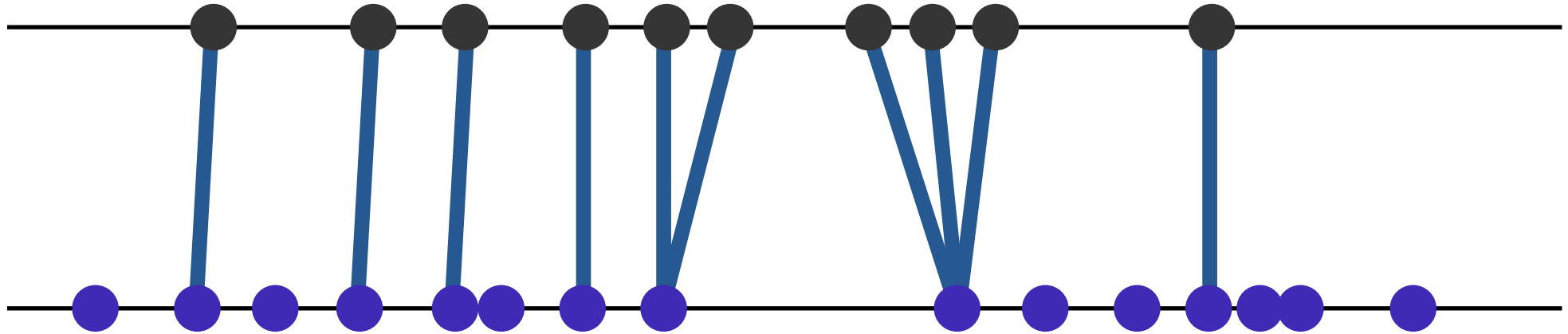
=> Sliced Partial Optimal Transport, [Bonneel and Coeurjolly 2019]



$$W(f, g) = \min \sum_{i,j} c_{i,j} \pi_{i,j} \quad \text{s.t.} \quad \begin{aligned} \sum_j \pi_{i,j} &= 1 \\ \sum_i \pi_{i,j} &\leq 1 \\ \pi_{i,j} &\geq 0 \end{aligned}$$

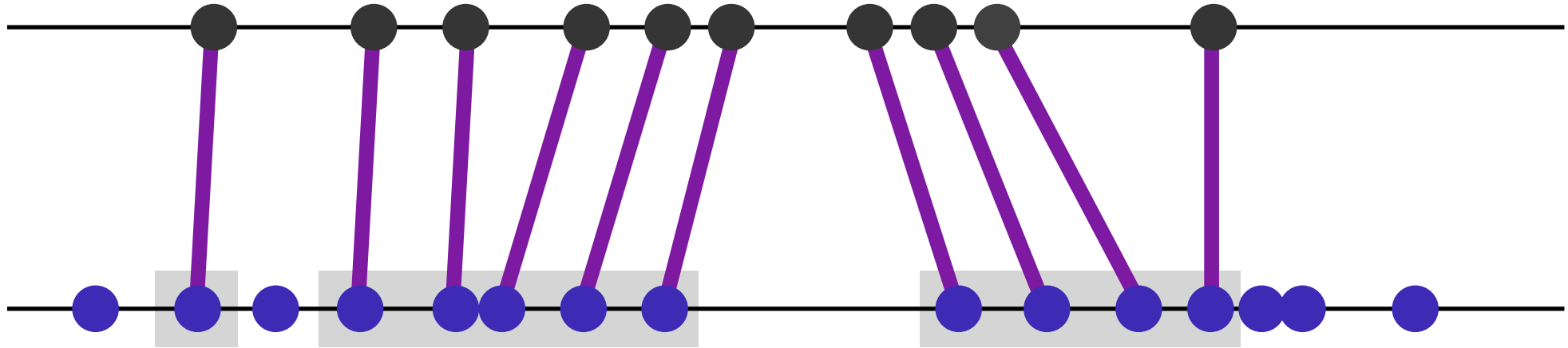
$$T \min_{\text{injective}} \sum_i c(x_i, y_{T(i)})$$

Quadratic time complexity algorithm (linear space)



— Euclidean Nearest Neighbor assignment

Quadratic time complexity algorithm (linear space)

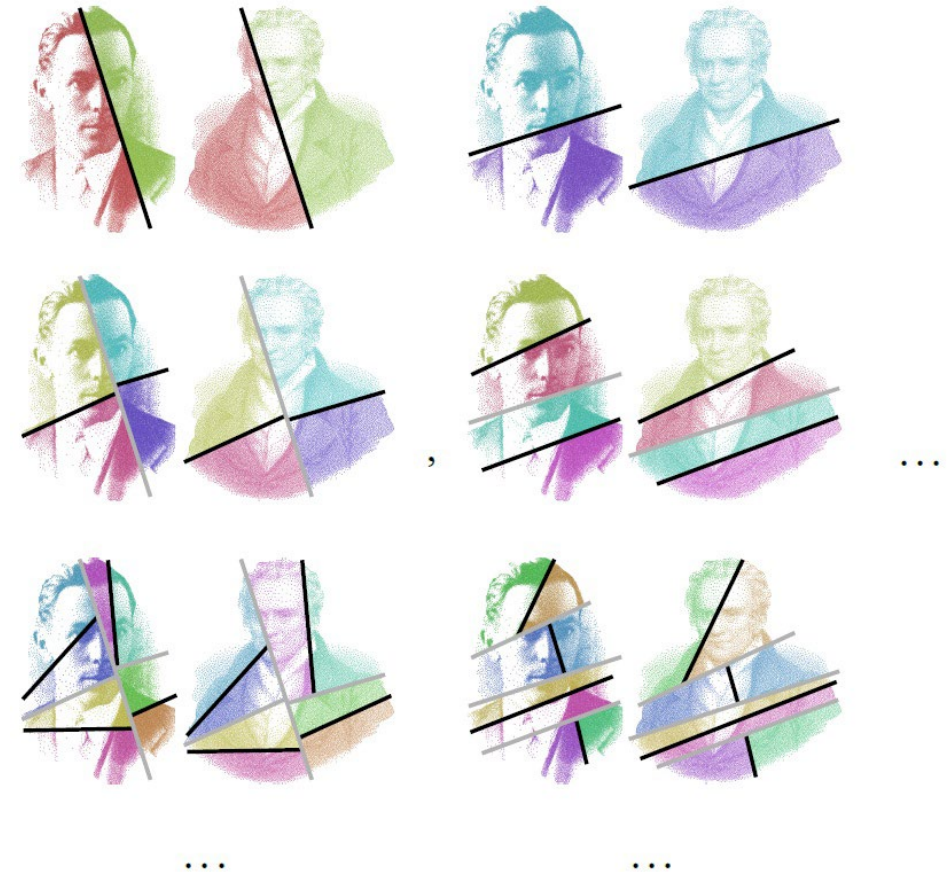


- Euclidean Nearest Neighbor assignment
- Optimal Transport assignment
- Intervals of bijective assignments

BSP OT

=> BSP-OT: Sparse transport plans between discrete measures in loglinear time, [Genest et al. 2025]

- Replaces the sorting step by a variant of QuickSort
- Each partition step of QuickSort splits the two input measures in two parts along one random direction
 - Each side of both point clouds has the same number of points
 - Results in a BSP tree





Semi-discrete optimal transport

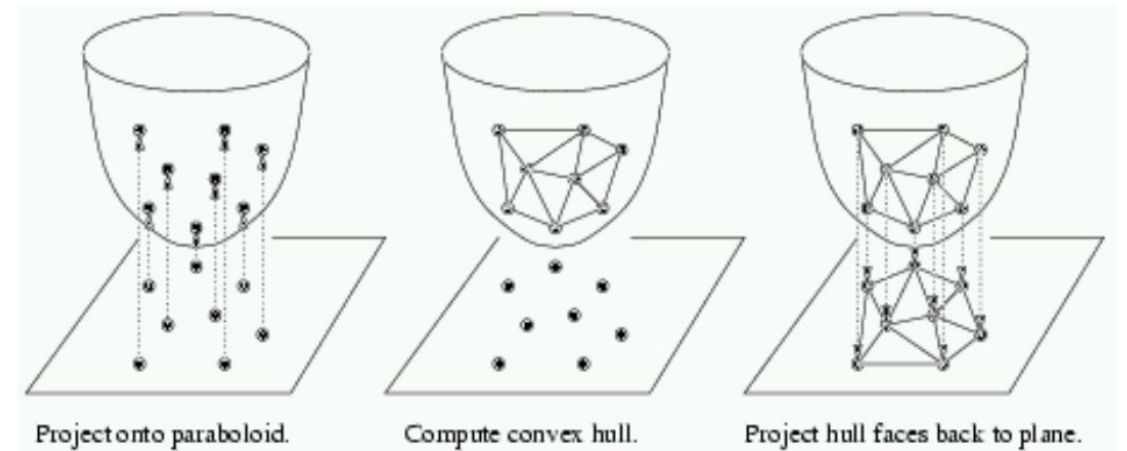
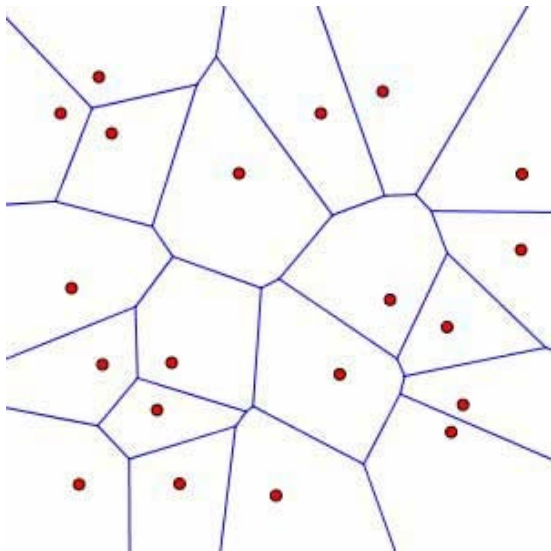
Voronoi diagram

- A partition such that each point x is assigned to its closest site x_i

$$\|x - x_i\|^2 \leq \|x - x_j\|^2 \quad \forall j$$

- The dual of a Delaunay triangulation: a triangulation of the sites such that no other site is encompassed by the circumcircle of a triangle

- Also: convex hull of a parabolic lifting



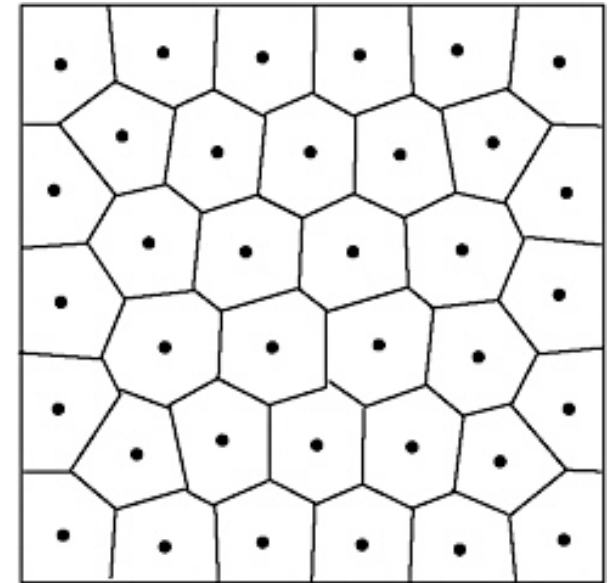
Centroidal Voronoi Diagram

- Can be defined as the solution to a least-square problem

$$\min \int_{Vor_i} \sum_i \|x - x_i\|^2 dx$$

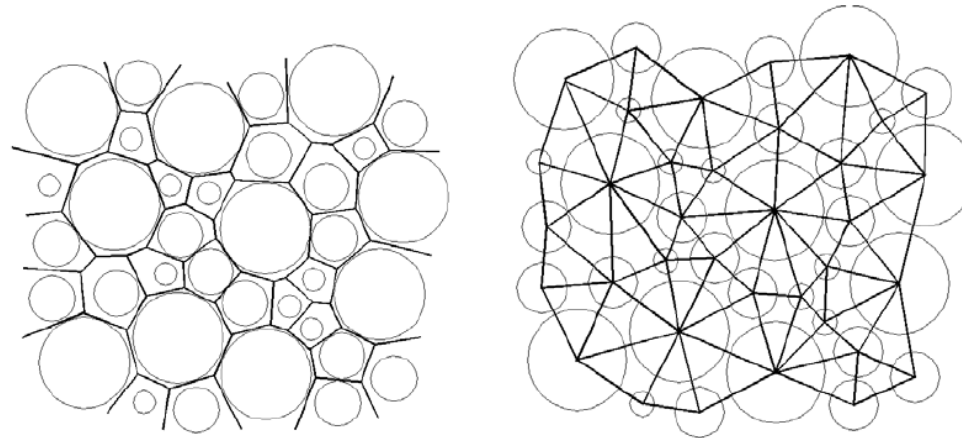
Also says that the centroid of Vor_i is the site x_i

- Can be computed by:
 - A Lloyd clustering algorithm
 - A descent approach on the above energy

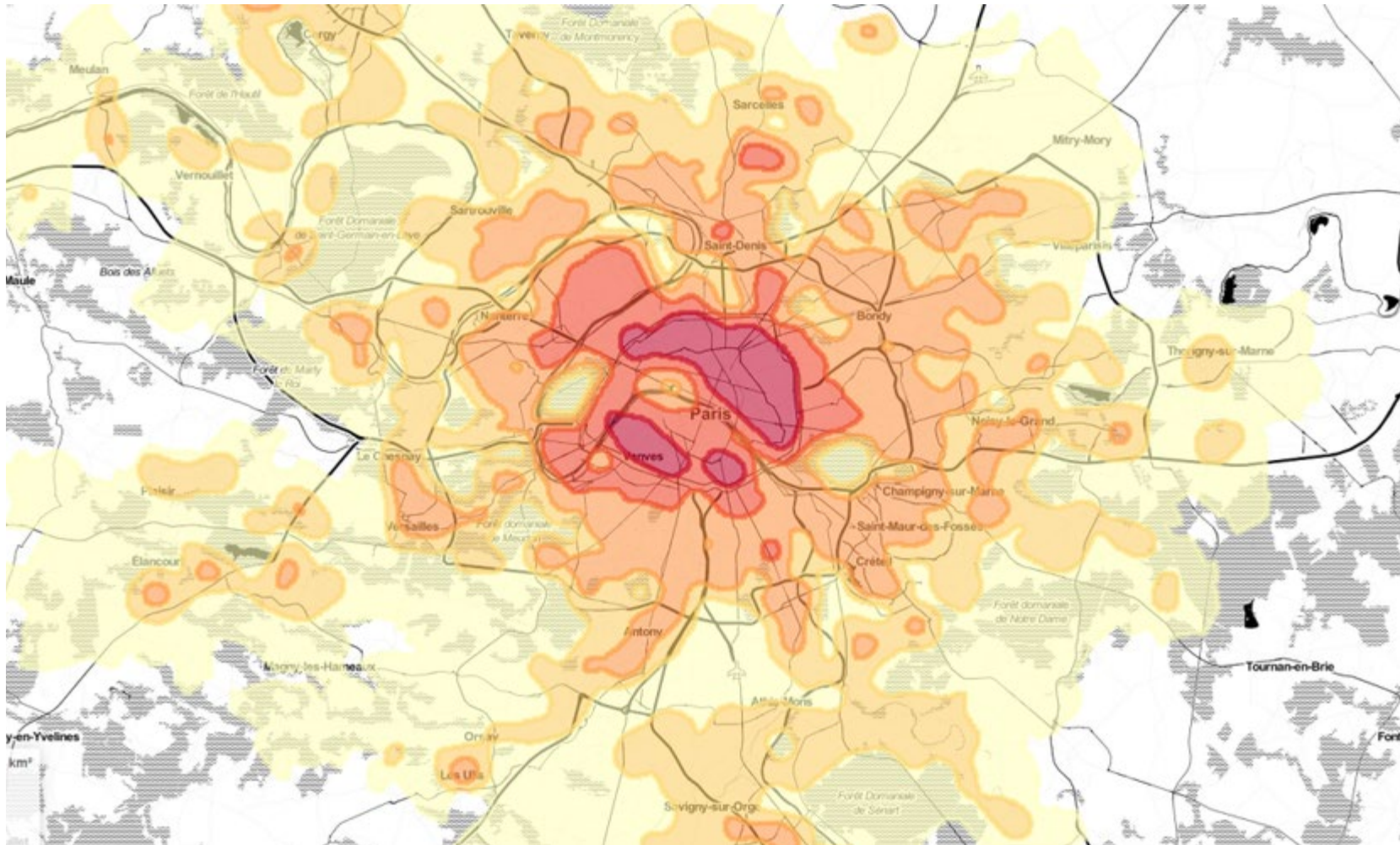


Power diagram (Laguerre diagram)

- A partition s.t. each point x is assigned to its closest site x_i with weight w_i
$$\|x - x_i\|^2 - w_i \leq \|x - x_j\|^2 - w_j \quad \forall j$$
- Can be computed by lifting a Voronoi diagram
 - Consider site coordinates $x'_i = (x_i; \sqrt{c - w_i})$ for large constant c ; $x' = (x; 0)$
 - Then $\|x' - x'_i\|^2 \leq \|x' - x'_j\|^2 \quad \forall j$
- Any partition into convex polyhedral cells is a power diagram of some sites

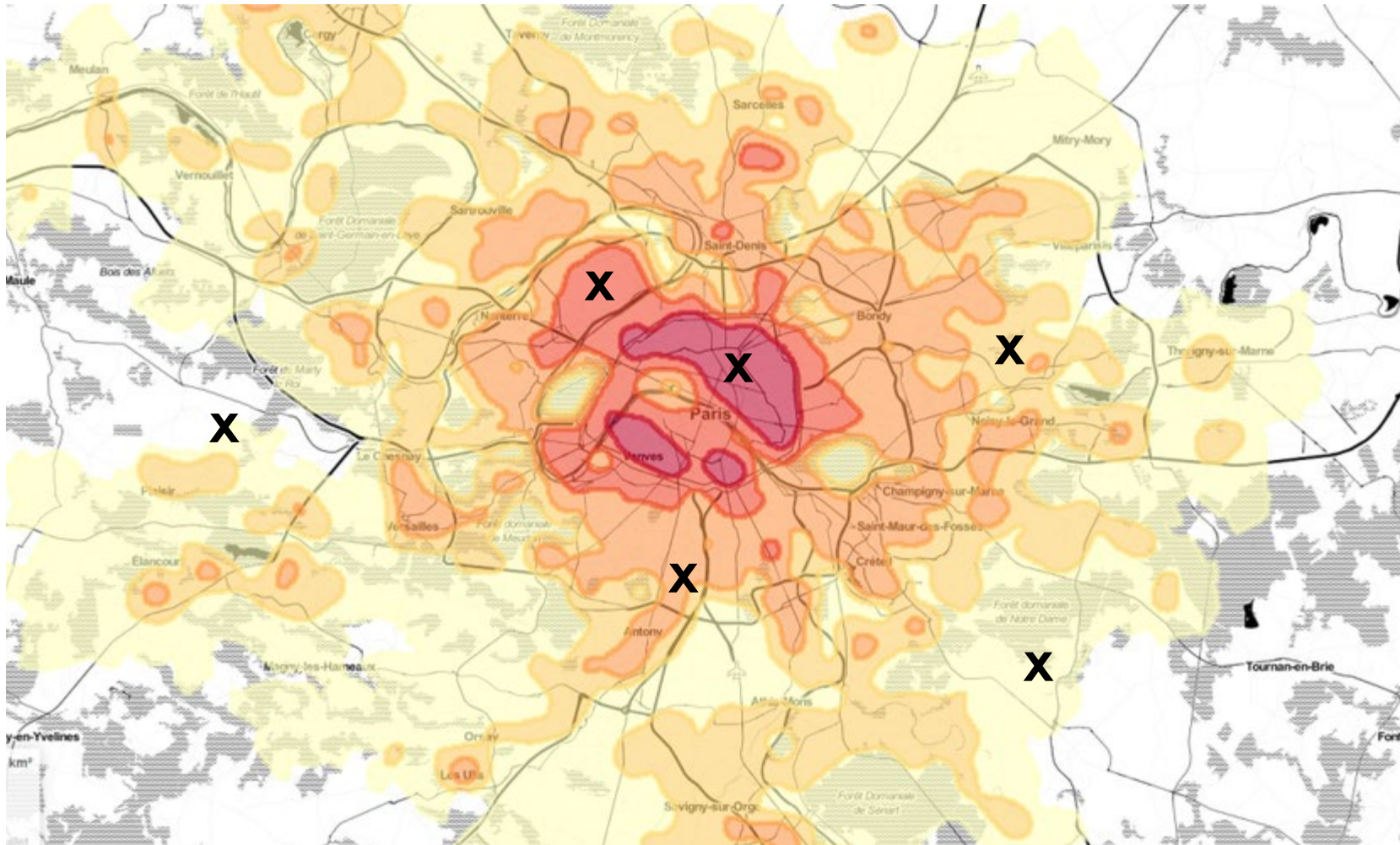


Semi-discrete Optimal Transport



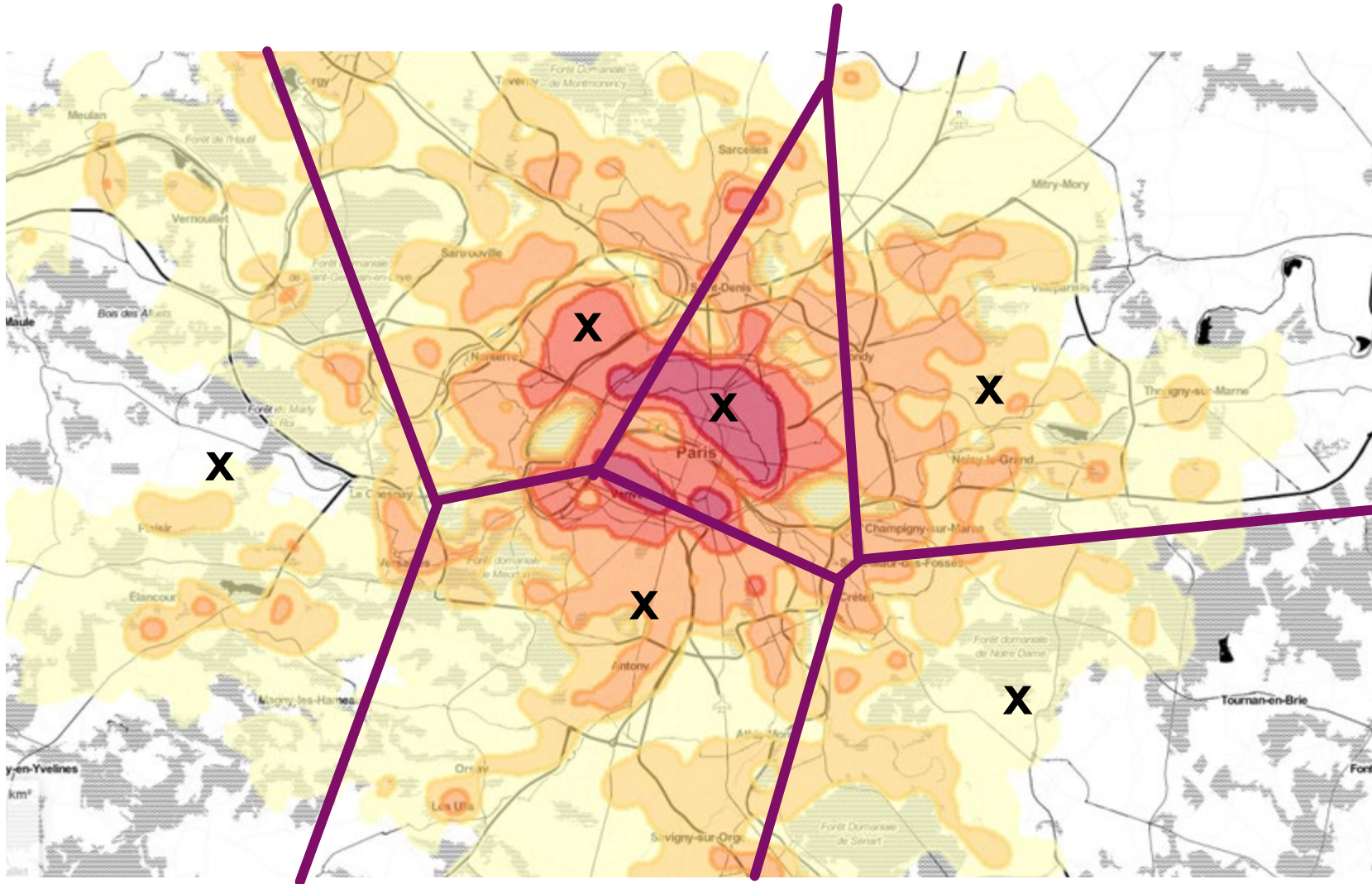
Population density f

Semi-discrete Optimal Transport



Set of bakeries, factories, ...?

Semi-discrete Optimal Transport



No constraint on production: population go to their nearest bakery/factory/... regardless of population

Semi-discrete Optimal Transport



Limited production: population go to the nearest bakery/factory **with sufficient production!**

Semi-discrete Optimal Transport



Limited production: population go to the nearest bakery/factory **with sufficient production!**



Back to optimal transport

► Optimal transport (Monge version) :

$$\min \int \|x - T(x)\|^2 d\mu(x)$$

Considering μ is continuous with density ρ

$$\min \int \|x - T(x)\|^2 \rho(x) dx$$

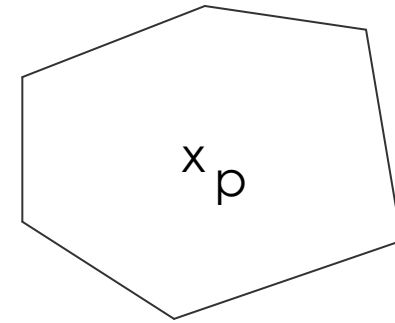
Considering ν (the target measure) discrete: $\nu = \sum \lambda_p \delta_p$

The mass preservation constraint is:

$$\lambda_p = \int_{T^{-1}(\{p\})} \rho(x) dx$$

Back to optimal transport

- In this case : $T^{-1}(\{p\}) = Vor^w(p)$
a power cell for some weight w_p

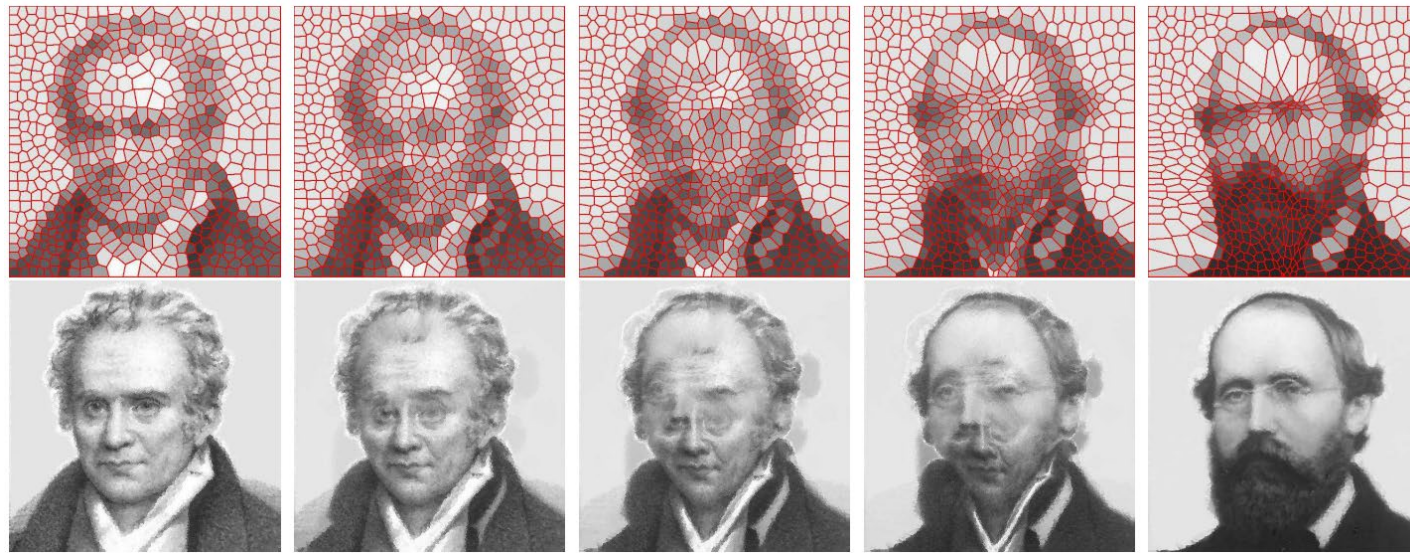


- This determines a partition, so Monge problem is:

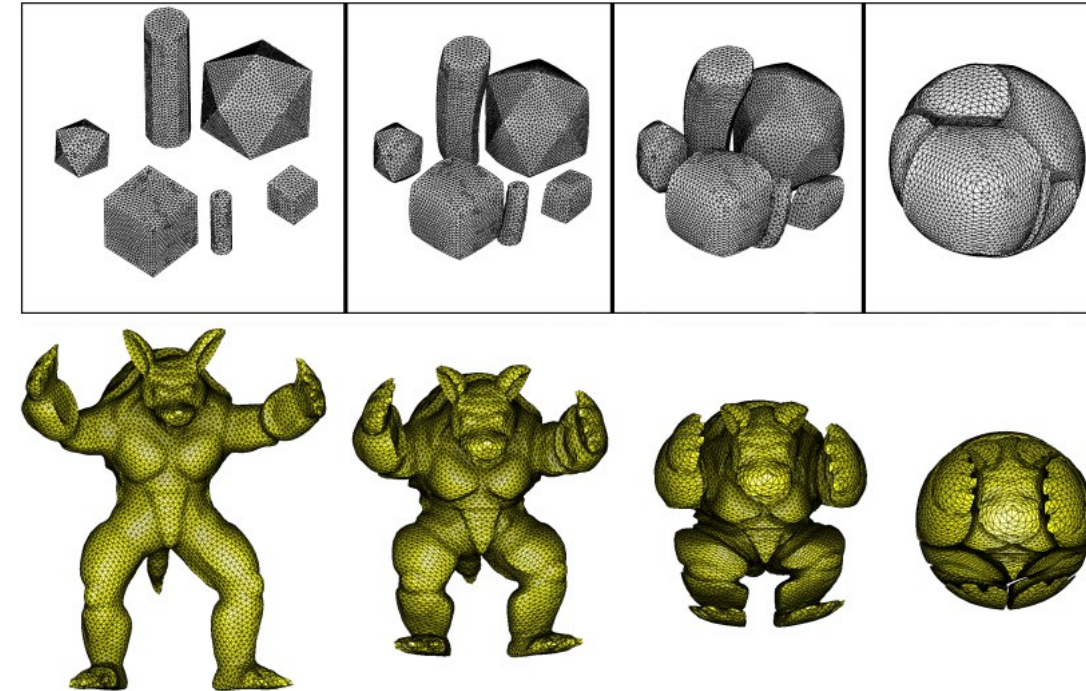
$$\min \sum_p \int_{Vor^w(p)} \|x - p\|^2 \rho(x) dx$$

- Idea: optimize weights w for each site to grow/shrink power cells until $\lambda_p = \int_{T^{-1}(\{p\})} \rho(x) dx$
- Gradient of appropriate functional given by $\frac{\partial \phi}{\partial w(p)}(w) = \lambda_p - \int_{Vor^w(p)} \rho(x) dx$

Back to optimal transport



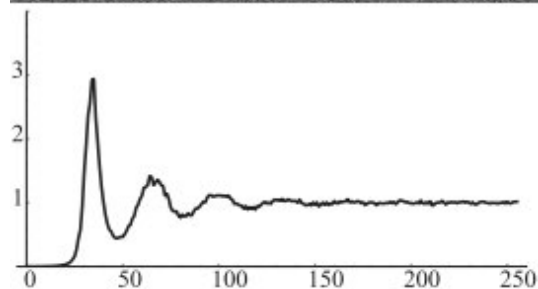
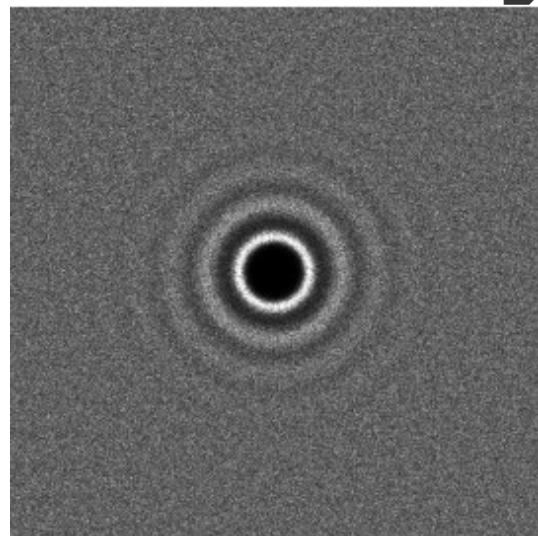
A Multiscale Approach to Optimal Transport [Mérigot 2011]



A Numerical Algorithm for L2 Semi-discrete Optimal Transport in 3D [Lévy 2015]

Application

Also optimizes for the locations p



Blue Noise through Optimal Transport [de Goes et al. 2012]